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Research article

(2n-3)-fault-tolerant Hamiltonian connectivity of augmented cubes AQ_n

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Abstract: The augmented cube AQ_n is an outstanding variation of the hypercube Q_n . It possesses many of the favorable properties of Q_n as well as some embedded properties not found in Q_n . This paper focuses on the fault-tolerant Hamiltonian connectivity of AQ_n . Under the assumption that $F \subset$ $V(AQ_n) \cup E(AQ_n)$ with $|F| \leq 2n - 3$, we proved that for any two different correct vertices u and v in AQ_n , there exists a fault-free Hamiltonian path that joins vertices u and v with the exception of (u, v), which is a weak vertex-pair in $AQ_n - F(n \geq 4)$. It is worth noting that in this paper we also proved that if there is a weak vertex-pair in $AQ_n - F$, there is at most one pair. This paper improved the current result that AQ_n is 2n - 4 fault-tolerant Hamiltonian connected. Our result is optimal and sharp under the condition of no restriction to each vertex.

Keywords: network topological structure; graph theory; augmented cube AQ_n ; computer network reliability; fault tolerance; Hamiltonian connectivity **Mathematics Subject Classification:** 05C38, 68M15

1. Introduction

Interconnection networks are of particular interest in the study of parallel computing systems. An interconnection network can be modeled as a graph G = (V, E), where V expresses the vertex set and E expresses the edge set. Exploring the structures of G is essential for designing a suitable topology for such an interconnection network. Topology structure has a crucial impact on the overall performance of the network. It determines the physical distribution of network nodes/links and the connection relationship between them. It also determines the number of hops in message transmission and the length of links per hop. Therefore, the topology has a great influence on the delay and power consumption. In addition, because the topology determines the number of available transmission paths between nodes, it also affects the distribution of network traffic, network bandwidth and transmission performance.

The hypercube Q_n is one of the most prevalent interconnection networks among all the popular parallel network topologies that possess properties such as logarithmic diameter, high symmetry, linear bisection width. The *n*-dimensional augmented cube AQ_n is an outstanding variation of the hypercube Q_n . It possesses many of the favorable properties of Q_n as well as some embedded properties not found in Q_n . Much research has been conducted on this type of augmented cube, and it appears frequently in the literatures [1,3,4,12–14,16,20,22].

In the last several years, the path embedding problem has become one of the most-studied graph embedding problems, appearing prolifically in the literatures [5, 6, 10, 14]. The fault-tolerant path embedding problem has also been the subject of frequent investigation, as is evident in the literatures [3, 4, 7, 10, 11, 22-24].

A path (or cycle) is considered a Hamiltonian path (or Hamiltonian cycle) if it passes through every vertex of graph G once and only once. A graph is said to be Hamiltonian if it contains at least one Hamiltonian cycle. One of the most challenging contemporary problems in graph theory is identifying a necessary and sufficient condition for a graph to be considered Hamiltonian. A graph G is regarded as Hamiltonian connected if for any pair of distinct vertices u and v, a Hamiltonian path P_{uv} exists. Hamiltonian paths and cycles have applicability for practical problems such as online optimization of complex, flexible manufacturing systems [25], wormhole routing [26], deadlock-free routing and broadcasting algorithms [27]. Applicability for such problems have been a driving force for the study of networks embedded with Hamiltonian paths.

In large interconnection networks, vertices(or edges) show a propensity for faultiness. This faultiness demands attention, as fault-tolerance serves as an important index of a network's stability. A graph *G* can be said to be *k*-fault-tolerant Hamiltonian connected if G - F remains Hamiltonian connected for any $F \subset V(G) \cup E(G)$ with $|F| \le k$.

Hsu et al. [3] proved that $AQ_n (n \ge 1)$ is Hamiltonian-connected. They also showed that AQ_n is (2n - 3)-fault-tolerant Hamiltonian and (2n - 4)-fault-tolerant Hamiltonian connected for $n \ge 4$ even when faulty elements occur. Soon after, Wang et al. [4] proved that AQ_n is (2n - 5)-fault-tolerant Panconnected for $n \ge 3$. We improved this result and showed that if $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \le 2n - 4$, then for any two distinct error-free vertices u and v with distance d, there exists an error-free path P_{uv} of length l with $max\{d + 2, 4\} \le l \le 2^n - f_v - 1$ in $AQ_n - F(n \ge 4)$ [22].

In this paper, we show that for any two distinct error-free vertices u and v, there exists a fault-free Hamiltonian path P_{uv} that joins vertices u and v with the exception of (u, v), which is a weak vertex-pair in $AQ_n - F(n \ge 4)$ under the assumption that $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \le 2n - 3$.

The rest of this paper is outlined as follows. Section 2 introduces the definitions and properties of the augmented cubes AQ_n . In Section 3, we investigate some lemmas of AQ_n to be used in our proofs. Section 4 proves the main theorem. Finally, Section 5 concludes the paper.

2. The definition and properties of AQ_n

In this section, we will introduce the definition of AQ_n and basic properties used in this paper.

Definition 2.1. The *n*-dimensional augmented cube AQ_n , proposed by Choudum and Sunitha [16], can be defined recursively as follows.

 AQ_1 is a complete graph K_2 with vertex set {0, 1}. For $n \ge 2$, AQ_n is obtained by taking two copies of the augmented cube AQ_{n-1} denoted by ${}^{0}AQ_{n-1}$ and ${}^{1}AQ_{n-1}$, and adding $2 \times 2^{n-1}$ edges between the

two copies as follows.

Let $V({}^{0}AQ_{n-1}) = \{0x_{n-1} \dots x_2x_1 | x_i \in \{0, 1\}\}$ and $V({}^{1}AQ_{n-1}) = \{1x_{n-1} \dots x_2x_1 | x_i \in \{0, 1\}\}$. A vertex $x \in V({}^{0}AQ_{n-1})$ is adjacent to a vertex $y \in V({}^{1}AQ_{n-1})$ if and only if either:

(1) $x_i = y_i$ for $1 \le i \le n - 1$; in this case, xy is called a hypercube edge and we set $x = y^{h_n}$ or $y = x^{h_n}$, or

(2) $x_i = \overline{y_i}$ for $1 \le i \le n - 1$; in this case, xy is called a complement edge and we set $x = y^{c_n}$ or $y = x^{c_n}$.

And an edge between $x = x_n x_{n-1} \dots x_i \dots x_2 x_1$ and $y = x_n x_{n-1} \dots \overline{x_i} \dots x_2 x_1 (x_i \in \{0, 1\}, 2 \le i \le n)$ is called an *i*-dimensional hypercube edge, setting $x = y^{h_i}$ or $y = x^{h_i}$, an edge between $x = x_n x_{n-1} \dots x_i \dots x_2 x_1$ and $y = x_n x_{n-1} \dots \overline{x_i} \dots \overline{x_2 x_1} (x_i \in \{0, 1\}, 1 \le i \le n)$ is called an *i*-dimensional complement edge, setting $x = y^{c_i}$ or $y = x^{c_i}$. For any vertex u in AQ_n , we use u^h to denote u^{h_n} and use u^c to denote u^{c_n} . Examples of the augmented cubes AQ_1 , AQ_2 and AQ_3 are shown in Figure 1(*a*)–(*c*), respectively.



 AQ_n is (2n - 1)-regular that has, naturally, 2^n vertices. For the sake of simplicity, we write this as $L = {}^0AQ_{n-1}$ and $R = {}^1AQ_{n-1}$. Let $N_{AQ_n}(x)$ express the set of vertices which are incident to vertex x and $E_{AQ_n}(x)$ express the set of edges which are incident to vertex x. For each vertex $x \in V(L)$ (or V(R)), let $N_L(x)$ (or $N_R(x)$) express the set of vertices adjacent to x in V(L)(or V(R)), E^C signify the set of edges joining L to R and $E_L(x)$ (or $E_R(x)$) serve as the set of edges which are incident to vertex x in L(or R).

For two distinct vertices u and v in G, a path P_{uv} from vertex u to vertex v of length k is a sequence of different vertices (x_0, x_1, \ldots, x_k) , where $x_0 = u$, $x_k = v$, $x_{i-1}x_i \in E(G)$ for each $i = 1, 2, \ldots, k$, and kis the number of edges in P_{uv} , called the length of P_{uv} . The distance between vertices u and v, denoted by d_{uv} , is the length of a shortest path from vertex u to vertex v in G.

Let $P_{uv} = (u, u_1, u_2, \dots, u_{k-1}, v)$ be a path from vertex u to vertex v. A subsequence $(u_i, u_{i+1}, \dots, u_j)$ of P_{uv} is called a subpath of P_{uv} , denoted by $P_{uv}(u_i, u_j)$. We can write $P_{uv} = (u, u_1, \dots, u_i, P_{uv}(u_i, u_j), u_j, \dots u_{k-1}, v)$. A cycle is a path P_{uv} with u = v. We use $C = (u, u_1, u_2, \dots, u_{k-1}, u)$ to denote a cycle containing vertex u.

We first give the useful definitions and properties about AQ_n as follows.

Proposition 2.1. Let uv be an edge in $AQ_n (n \ge 2)$ [4]. If uv is not an (n - 1)-dimensional complement edge, then u^h , u^c , v^h and v^c are all distinct. Otherwise, $u^h = v^c$, $u^c = v^h$.

Proposition 2.2. For any two distinct vertices *u* and *v* of AQ_n , we have $|(N_{AQ_n}(u) \cup E_{AQ_n}(u)) \cap (N_{AQ_n}(v) \cup E_{AQ_n}(v))| \le 5$.

Proof. Let *m* be the adjacent vertex of vertex *u*. Then

$$m = \begin{cases} u^{c_j} = u_n u_{n-1} \cdots u_{j+1} \bar{u}_j \bar{u}_{j-1} \cdots \bar{u}_2 \bar{u}_1, & 1 \le j \le n. \\ u^{h_j} = u_n u_{n-1} \cdots u_{j+1} \bar{u}_j u_{j-1} \cdots u_2 u_1, & 2 \le j \le n. \end{cases}$$

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If there exists at least one common adjacent vertex between any two distinct vertices u and v of AQ_n , then $d_{uv} \le 2$. We divided to the following two cases to prove.

Case 1. $d_{uv} = 1$.

Case 1.1. $v = u^{c_i} (1 \le i \le n)$. Then $v = u_n u_{n-1} \cdots u_{i+1} \overline{u}_i \overline{u}_{i-1} \cdots \overline{u}_2 \overline{u}_1$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{c_{i+1}} = u_n u_{n-1} \cdots \bar{u}_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, & 1 \le i \le n-1. \\ u^{c_{i-1}} = u_n u_{n-1} \cdots u_{i+1} u_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, & 2 \le i \le n. \\ u^{h_i} = u_n u_{n-1} \cdots u_{i+1} \bar{u}_i u_{i-1} \cdots u_2 u_1, & 2 \le i \le n. \\ u^{h_{i+1}} = u_n u_{n-1} \cdots \bar{u}_{i+1} u_i u_{i-1} \cdots u_2 u_1, & 1 \le i \le n-1. \end{cases}$$

Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_{i+1}}, u^{c_{i-1}}, u^{h_i}, u^{h_{i+1}}\}$ for $2 \le i \le n-1$, $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_{i+1}}, u^{h_{i+1}}\}$ for i = 1, $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_{i-1}}, u^{h_i}\}$ for i = n.

Case 1.2. $v = u^{h_i} (2 \le i \le n)$. Then $v = u_n u_{n-1} \cdots u_{i+1} \bar{u}_i u_{i-1} \cdots u_2 u_1$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{c_{i-1}} = u_n u_{n-1} \cdots u_{i+1} u_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, & 2 \le i \le n. \\ u^{c_i} = u_n u_{n-1} \cdots u_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, & 2 \le i \le n. \end{cases}$$

Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_{i-1}}, u^{c_i}\}$ for $2 \le i \le n$. **Case 2.** $d_{uv} = 2$. **Case 2.1** $v = u^{c_i c_k} (1 \le i \le k \le n, k \ne i \pm 1)$ Then $v = u, u = \dots u$

Case 2.1. $v = u^{c_i c_k} (1 \le i < k \le n, k \ne i + 1)$. Then $v = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k \bar{u}_{k-1} \cdots \bar{u}_{i+1} u_i u_{i-1} \cdots u_2 u_1$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{c_i} = u_n u_{n-1} \cdots u_{k+1} u_k u_{k-1} \cdots u_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \\ u^{c_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k \bar{u}_{k-1} \cdots \bar{u}_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \\ u^{h_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k u_{i+1} u_i u_{i-1} \cdots u_2 u_1, \qquad k = i+2. \\ u^{h_{i+1}} = u_n u_{n-1} \cdots u_{k+1} u_k \bar{u}_{i+1} u_i u_{i-1} \cdots u_2 u_1, \qquad k = i+2. \end{cases}$$

Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_i}, u^{c_k}, u^{h_k}, u^{h_{i+1}}\}$ for k = i + 2 and $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_i}, u^{c_k}\}$ for $k \neq i + 2$.

Case 2.2. $v = u^{c_i h_k} (1 \le i < k \le n, k \ne i + 1)$. Then $v = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k u_{k-1} \cdots u_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{c_i} = u_n u_{n-1} \cdots u_{k+1} u_k u_{k-1} \cdots u_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \\ u^{h_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k u_{k-1} \cdots u_{i+1} u_i u_{i-1} \cdots u_2 u_1, \\ u^{c_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k \bar{u}_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \qquad k = i+2. \\ u^{h_{i+1}} = u_n u_{n-1} \cdots u_{k+1} u_k \bar{u}_{i+1} u_i u_{i-1} \cdots u_2 u_1, \qquad k = i+2. \end{cases}$$

Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_i}, u^{h_k}, u^{c_k}, u^{h_{i+1}}\}$ for k = i + 2 and $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{c_i}, u^{h_k}\}$ for $k \neq i + 2$.

Case 2.3. $v = u^{h_i h_k} (2 \le i < k \le n)$. Then $v = u_n u_{n-1} \cdots u_{k+1} \overline{u}_k u_{k-1} \cdots u_{i+1} \overline{u}_i u_{i-1} \cdots u_{2} u_1$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{h_i} = u_n u_{n-1} \cdots u_{k+1} u_k u_{k-1} \cdots u_{i+1} \overline{u}_i u_{i-1} \cdots u_2 u_1, \\ u^{h_k} = u_n u_{n-1} \cdots u_{k+1} \overline{u}_k u_{k-1} \cdots u_{i+1} u_i u_{i-1} \cdots u_2 u_1, \\ u^{c_{i-1}} = u_n u_{n-1} \cdots u_{k+1} u_k u_i \overline{u}_{i-1} \cdots \overline{u}_2 \overline{u}_1, \qquad k = i+1. \\ u^{c_k} = u_n u_{n-1} \cdots u_{k+1} \overline{u}_k \overline{u}_i \overline{u}_{i-1} \cdots \overline{u}_2 \overline{u}_1, \qquad k = i+1. \end{cases}$$

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Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{h_i}, u^{h_k}, u^{c_{i-1}}, u^{c_k}\}$ for k = i + 1 and $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{h_i}, u^{h_k}\}$ for $k \neq i + 1$.

Case 2.4. $v = u^{h_i c_k} (2 \le i < k \le n)$. Then $v = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k \bar{u}_{k-1} \cdots \bar{u}_{i+1} u_i \bar{u}_{i-1} \cdots \bar{u}_{2\bar{u}_{1}}$. $m \in N(v)$ if and only if

$$m = \begin{cases} u^{h_i} = u_n u_{n-1} \cdots u_{k+1} u_k u_{k-1} \cdots u_{i+1} \bar{u}_i u_{i-1} \cdots u_2 u_1, \\ u^{c_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k \bar{u}_{k-1} \cdots \bar{u}_{i+1} \bar{u}_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \\ u^{c_{i-1}} = u_n u_{n-1} \cdots u_{k+1} u_k u_i \bar{u}_{i-1} \cdots \bar{u}_2 \bar{u}_1, \qquad k = i+1. \\ u^{h_k} = u_n u_{n-1} \cdots u_{k+1} \bar{u}_k u_i u_{i-1} \cdots u_2 u_1, \qquad k = i+1. \end{cases}$$

Then $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{h_i}, u^{c_k}, u^{c_{i-1}}, u^{h_k}\}$ for k = i + 1 and $N_{AQ_n}(u) \cap N_{AQ_n}(v) = \{u^{h_i}, u^{c_k}\}$ for $k \neq i + 1$.

Combining the above cases, we have that $|N_{AQ_n}(u) \cap N_{AQ_n}(v)| \le 4$. If $uv \in E(AQ_n)$, then $|(N_{AQ_n}(u) \cup E_{AQ_n}(u)) \cap (N_{AQ_n}(v) \cup E_{AQ_n}(v))| \le 5$. If $uv \notin E(AQ_n)$, then $|(N_{AQ_n}(u) \cup E_{AQ_n}(u)) \cap (N_{AQ_n}(v) \cup E_{AQ_n}(v))| \le 4$. Hence, we have $|(N_{AQ_n}(u) \cup E_{AQ_n}(u)) \cap (N_{AQ_n}(v) \cup E_{AQ_n}(v))| \le 5$.

Definition 2.2. Let $F \subset V(AQ_n) \cup E(AQ_n)$ with |F| = 2n - 3. If $AQ_n - F$ include a vertex w with $N_{AQ_n-F}(w) = \{w_1, w_2\}$, then w is considered a weak 2-degree vertex and (w_1, w_2) is considered a w-weak vertex pair(or a weak vertex pair, for short).

Take $F = \{a, b\}$ for an example, we know that w is a weak 2-degree vertex and (w_1, w_2) is a weak vertex-pair in $AQ_3 - F$ (See Figure 2).



Figure 2. Illustration of weak vertex-pair.

Because it is impossible for any error-free path $P_{w_1w_2}$ of length l with $l \ge 3$ to contain the weak 2-degree vertex w, no error-free Hamiltonian path that joins vertices w_1 and w_2 exists in $AQ_n - F$. Fortunately, at most one weak 2-degree vertex w and at most one w-weak vertex-pair exist in $AQ_n - F(n \ge 5)$ for any $F \subset V(AQ_n) \cup E(AQ_n)$ with |F| = 2n - 3. We shall provide proof of this fact in Proposition 2.3.

Definition 2.3. If (w_1, w_2) is not a weak vertex-pair for arbitrary vertex $w \in V(AQ_n - F)$, then (w_1, w_2) is called a normal vertex pair.

Proposition 2.3. Let $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \le 2n - 1 (n \ge 6)$. Then there exists at most one vertex $z \in V(AQ_n - F)$ such that $d_{AQ_n - F}(z) \le 2$.

Proof. Assume that there exist two vertices z_1, z_2 such that $d_{AQ_n-F}(z_1) \le 2$ and $d_{AQ_n-F}(z_2) \le 2$. Since AQ_n is a (2n - 1)-regular graph, we have $|F \cap (N_{AQ_n}(z_1) \cup E_{AQ_n}(z_1))| \ge 2n - 3$ and $|F \cap (N_{AQ_n}(z_2) \cup E_{AQ_n}(z_2))| \ge 2n - 3$.

By Proposition 2.2, $|(N_{AQ_n}(z_1) \cup E_{AQ_n}(z_1)) \cap (N_{AQ_n}(z_2) \cup E_{AQ_n}(z_2))| \le 5$. Then $|F \cap (N_{AQ_n}(z_1) \cup E_{AQ_n}(z_1))| + |F \cap (N_{AQ_n}(z_2) \cup E_{AQ_n}(z_2))| - |F \cap (N_{AQ_n}(z_1) \cup E_{AQ_n}(z_1)) \cap (N_{AQ_n}(z_2) \cup E_{AQ_n}(z_2))| \ge 2(2n-3)-5 = 4n-11 > 2n-1 (n \ge 6)$, a contradiction to $|F| \le 2n-1$.

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Hence, at most one vertex $z \in V(AQ_n - F)$ exists in $AQ_n - F$ such that $d_{AQ_n - F}(z) \leq 2$ for any $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \leq 2n - 1$.

By Proposition 2.3, we can obtain the Corollary 2.1 as follows.

Corollary 2.1. Let $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \le 2n - 3$. Then at most a weak 2-degree vertex *w* and at most a *w*-weak vertex-pair exist in $AQ_n - F(n \ge 6)$.

3. Some lemmas

Denote $F^{L} = F \cap L$, $F^{R} = F \cap R$, $F^{C} = F \cap E^{C}$, $F_{v} = F \cap V(AQ_{n})$, $F_{e} = F \cap E(AQ_{n})$, $F_{v}^{L} = V(L) \cap F^{L}$, $F_{v}^{R} = V(R) \cap F^{R}$, $f_{v} = |F_{v}|$, $f_{v}^{L} = |F_{v}^{L}|$, $f_{v}^{R} = |F_{v}^{R}|$.

We need the following lemmas.

Lemma 3.1. $AQ_n (n \ge 3)$ is (2n - 4)-fault-tolerant Hamiltonian connected and (2n - 3)-fault-tolerant Hamiltonian [3].

Lemma 3.2. If $|F^L| = 2n - 3(n \ge 6)$, then for arbitrary vertex-pair (u, v) with $u, v \in V(L)$, there exist two faulty elements $x_1, x_2 \in F^L$ such that (u, v) is a normal vertex-pair in $L - (F^L - \{x_1, x_2\})$.

Proof. We use the following two cases to prove.

Case 1. (u, v) is a *w*-weak vertex pair in $L - F^L$. Then $d_{L-F^L}(w) = 2$ and $|F^L \cap (N_L(w) \cup E_L(w))| = 2n - 5$. By Proposition 2.3, *w* is the unique vertex in $L - F^L$ with $d_{L-F^L}(w) \le 2$.

So, we can choose two elements $x_1, x_2 \in F^L \cap (N_L(w) \cup E_L(w))$. Let $F_1^L = F^L - \{x_1, x_2\}$, then $d_{L-F_1^L}(w) = 4$. Hence, (u, v) is a normal vertex pair in $L - F_1^L$.

Case 2. (u, v) is a normal vertex pair in $L - F^L$. Let vertex $z \in L - F^L$ with the least degree.

Case 2.1. $\delta(L - F^L) = 0$, then $F^L \subset N_L(z) \cup E_L(z)$, $d_{L-F^L}(z) = 0$.

Case 2.1.1. $z \in \{u, v\}$. Assume that u = z (When v = z, the same proofs apply).

By Proposition 2.3, *u* is the unique vertex in $L - F^L$ with $d_{L-F^L}(u) \le 2$. Then (u, v) is a normal vertex-pair in $L - (F^L - \{x_1, x_2\})$ for arbitrary two elements x_1, x_2 of F^L .

Case 2.1.2. $z \notin \{u, v\}$.

By Proposition 2.3, z is the unique vertex in $L - F^L$ with $d_{L-F^L}(z) \le 2$. There exist two elements $x_1, x_2 \in F^L \cap (N_L(z) \cup E_L(z))$ with $x_1, x_2 \notin \{uz, vz\}$. Let $F_1^L = F^L - \{x_1, x_2\}$, then $d_{L-F_1^L}(z) = 2$ and $N_{L-F_1^L}(z) \neq \{u, v\}$. Hence, (u, v) is a normal vertex pair in $L - F_1^L$.

Case 2.2. $\delta(L - F^L) = 1$. Then $|F^L \cap (N_L(z) \cup E_L(z))| = 2n - 4$ and $d_{L-F^L}(z) = 1$.

By Proposition 2.3, z is the unique vertex in $L - F^L$ with $d_{L-F^L}(z) \le 2$. There exist two elements $x_1, x_2 \in F^L \cap (N_L(z) \cup E_L(z))$. Let $F_1^L = F^L - \{x_1, x_2\}$, then $d_{L-F_1^L}(z) = 3$. Then (u, v) is a normal vertex pair in $L - F_1^L$.

Case 2.3. $\delta(L - F^L) \ge 2$.

Then (u, v) is a normal vertex-pair in $L - (F^L - \{x_1, x_2\})$ for arbitrary two elements x_1, x_2 of F^L . The Lemma holds.

By a similar argument to Lemma 3.2, we can get the following Lemma.

Lemma 3.3. If $|F^L| = 2n - 4(n \ge 6)$, then for arbitrary vertex-pair (u, v) with $u, v \in V(L)$, there is a faulty element $x \in F^L$ such that (u, v) is a normal vertex-pair in $L - (F^L - \{x\})$.

Proof. We prove the lemma by induction on $n \ge 4$. The induction basis for n = 4 holds by computer program. Suppose that the lemma holds for n - 1 with $n \ge 5$, then we must show the lemma holds for n.

Case 1. $x, y, z, w \in V(L)$.

By induction hypothesis, two disjoint paths P_{xy} , P_{zw} with $V(P_{xy}) \cup V(P_{zw}) = V(L - F^L)$ exist in $L - F^L$. Assume that $|V(P_{xy})| \ge |V(P_{zw})|$. Then $\frac{|V(P_{xy})|}{2} \ge \lceil \frac{2^n - 2}{2} \rceil/2 \ge 7 (n \ge 5)$. Since $|F| \le 2$, we can choose an edge $x_1y_1 \subset E(P_{xy})$ such that $x_1^h, y_1^h, x_1x_1^h, y_1y_1^h \notin F$. By Lemma 3.1, a Hamiltonian path $P_{x_1^hy_1^h}$ will exist in $R - F^R$. Let $P_{xy}^1 = (x, P_{xy}(x, x_1), x_1, x_1^h, P_{x_1^hy_1^h}, y_1^h, y_1, P_{xy}(y_1, y), y)$. Then P_{xy}^1, P_{zw} are two desired paths in $AQ_n - F$ (see Figure 3(a)).

Case 2. $x, y, z \in V(L), w \in V(R)$.

Since $|E^C| - 12 = 2^n - 12 \ge 20 (n \ge 5)$, an error-free edge $z_1 z_1^h$ can be chosen form E^C such that $z_1 \notin \{x, y, z\}, z_1^h \neq w$ and $z_1, z_1^h \notin F$. By induction hypothesis, two disjoint paths P_{xy}, P_{zz_1} with $V(P_{xy}) \cup V(P_{zz_1}) = V(L - F^L)$ exist in $L - F^L$. By Lemma 3.1, we can find a Hamiltonian path $P_{z_1^h w}$ in $R - F^R$. Let $P_{zw} = (z, P_{zz_1}, z_1, z_1^h, P_{z_1^h w}, w)$. Then P_{xy}, P_{zw} are two desired paths in $AQ_n - F$ (see Figure 3(b)).



Figure 3. Illustrations of Lemma 3.4.

Case 3. $x, y \in V(L), z, w \in V(R)$.

By Lemma 3.1, we can find a Hamiltonian path P_{xy} in $L - F^L$ and a Hamiltonian path P_{zw} in $R - F^R$. Then P_{xy} , P_{zw} are two desired paths in $AQ_n - F$ (see Figure 4(a)).

Case 4. $x, z \in V(L), y, w \in V(R)$.

Since $|E^C| - 12 = 2^n - 12 \ge 20 (n \ge 5)$, there are two error-free disjoint edges $x_1 x_1^h, z_1 z_1^h$ such that $x_1, z_1 \notin \{x, z\}, x_1^h, z_1^h \notin \{y, w\}$ and $x_1, x_1^h, z_1, z_1^h \notin F$. By induction hypothesis, two disjoint paths P_{xx_1}, P_{zz_1} with $V(P_{xx_1}) \cup V(P_{zz_1}) = V(L - F^L)$ exist in $L - F^L$ and two disjoint paths $P_{x_1^h y}, P_{z_1^h w}$ with $V(P_{x_1^h y}) \cup V(P_{z_1^h w}) = V(R - F^R)$ exist in $R - F^R$. Let $P_{xy} = (x, P_{xx_1}, x_1, x_1^h, P_{x_1^h y}, y), P_{zw} = (z, P_{zz_1}, z_1, z_1^h, P_{z_1^h w}, w)$. Then P_{xy}, P_{zw} are two desired paths in $AQ_n - F$ (see Figure 4(b)).



Figure 4. Illustrations of Lemma 3.4.

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The lemma holds.

Lemma 3.5. Let x, y, z, w, a, b be six distinct vertices in $AQ_n (n \ge 4)$. Then there exist three disjoint paths P_{xy} , P_{zw} , and P_{ab} such that $V(P_{xy}) \cup V(P_{zw}) \cup V(P_{ab}) = V(AQ_n)$.

Proof. We prove the lemma by induction on $n \ge 4$. The induction basis for n = 4 holds by computer program. Suppose the lemma holds for n - 1 with $n \ge 5$, then we must show the lemma holds for n.

Case 1. $x, y, z, w, a, b \in V(L)$.

By induction hypothesis, three disjoint paths P_{xy} , P_{zw} and P_{ab} with $V(P_{xy}) \cup V(P_{zw}) \cup V(P_{ab}) = V(L)$ will exist in *L*. Let $x_1y_1 \in E(P_{xy})$. By Lemma 3.1, a Hamiltonian path $P_{x_1^h y_1^h}$ can be found in *R*. Let $P_{xy}^1 = (x, P_{xy}(x, x_1), x_1, x_1^h, P_{x_1^h y_1^h}, y_1^h, y_1, P_{xy}(y_1, y), y)$. Then P_{xy}^1, P_{zw}, P_{ab} are three desired paths in AQ_n (see Figure 5(a)).

Case 2. $x, y, z, w, a \in V(L), b \in V(R)$.

Choose a vertex a_1 from $V(L) - \{x, y, z, w, a\}$ such that there exists a vertex $a_1^g \in \{a_1^h, a_1^c\}$ with $a_1^g \neq b$. By induction hypothesis, three disjoint paths P_{xy} , P_{zw} and P_{aa_1} with $V(P_{xy}) \cup V(P_{zw}) \cup V(P_{aa_1}) = V(L)$ will exist in *L*. By Lemma 3.1, a Hamiltonian path $P_{a_1^g b}$ can be found in *R*. Let $P_{ab} = (a, P_{aa_1}, a_1, a_1^g, P_{a_1^g b}, b)$. Then P_{xy}, P_{zw}, P_{ab} are three desired paths in AQ_n (see Figure 5(b)).

Case 3. $x, y, z, w \in V(L), a, b \in V(R)$.

By Lemma 3.4, two disjoint paths P_{xy} , P_{zw} with $V(P_{xy}) \cup V(P_{zw}) = V(L)$ exist in L. By Lemma 3.1, we can find a Hamiltonian path P_{ab} in R. Then P_{xy} , P_{zw} , P_{ab} are three desired paths in AQ_n (see Figure 5(c)).



Figure 5. Illustrations of Lemma 3.5.

Case 4. $x, y, z, a \in V(L), w, b \in V(R)$.

Since $|E^{C}|-12 = 2^{n}-12 \ge 20$, there are two disjoint edges $z_{1}z_{1}^{h}, a_{1}a_{1}^{h}$ such that $z_{1}, a_{1} \notin \{x, y, z, a\}$ and $z_{1}^{h}, a_{1}^{h} \notin \{w, b\}$. By induction hypothesis, three disjoint paths $P_{xy}, P_{zz_{1}}$ and $P_{aa_{1}}$ with $V(P_{xy}) \cup V(P_{zz_{1}}) \cup V(P_{aa_{1}}) = V(L)$ will exist in *L*. And by Lemma 3.4, two disjoint paths $P_{z_{1}^{h}w}, P_{a_{1}^{h}b}$ with $V(P_{z_{1}^{h}w}) \cup V(P_{aa_{1}}) = V(R)$ will exist in *R*. Let $P_{zw} = (z, P_{zz_{1}}, z_{1}, z_{1}^{h}, P_{z_{1}^{h}w}, w)$ and $P_{ab} = (a, P_{aa_{1}}, a_{1}, a_{1}^{h}, P_{a_{1}^{h}b}, b)$. Then P_{xy}, P_{zw}, P_{ab} are three desired paths in AQ_{n} (see Figure 6(a)).

Case 5. $x, z, a \in V(L), y, w, b \in V(R)$.

Since $|E^{C}| - 12 = 2^{n} - 12 \ge 20$, there are three disjoint edges $x_{1}x_{1}^{h}, z_{1}z_{1}^{h}, a_{1}a_{1}^{h}$ such that $x_{1}, z_{1}, a_{1} \notin \{x, z, a\}$ and $x_{1}^{h}, z_{1}^{h}, a_{1}^{h} \notin \{y, w, b\}$. By induction hypothesis, three disjoint paths $P_{xx_{1}}, P_{zz_{1}}$ and $P_{aa_{1}}$ with $V(P_{xx_{1}}) \cup V(P_{zz_{1}}) \cup V(P_{aa_{1}}) = V(L)$ will exist in L and three disjoint paths $P_{x_{1}^{h}y}, P_{z_{1}^{h}w}, P_{a_{1}^{h}b}$ with $V(P_{x_{1}^{h}y}) \cup V(P_{z_{1}^{h}w}) \cup V(P_{a_{1}^{h}b}) = V(R)$ will exist in R. Let $P_{xy} = (x, P_{xx_{1}}, x_{1}, x_{1}^{h}, P_{x_{1}^{h}y}, y), P_{zw} = (z, P_{zz_{1}}, z_{1}, z_{1}^{h}, P_{z_{1}^{h}w}, w)$ and $P_{ab} = (a, P_{aa_{1}}, a_{1}, a_{1}^{h}, P_{a_{1}^{h}b}, b)$. Then P_{xy}, P_{zw}, P_{ab} are three desired paths in AQ_{n} (see Figure 6(b)).

Case 6. $x, y, a \in V(L), z, w, b \in V(R)$.

Since $|E^{C}| - 12 = 2^{n} - 12 \ge 20$, there is an edge $a_{1}a_{1}^{h}$ such that $a_{1} \notin \{x, y, a\}$ and $a_{1}^{h} \notin \{z, w, b\}$. By Lemma 3.4, two disjoint paths P_{xy} and $P_{aa_{1}}$ with $V(P_{xy}) \cup V(P_{aa_{1}}) = V(L)$ will exist in *L* and two disjoint paths P_{zw} and $P_{a_{1}^{h}b}$ with $V(P_{zw}) \cup V(P_{a_{1}^{h}b}) = V(R)$ will exist in *R*. Let $P_{ab} = (a, P_{aa_{1}}, a_{1}, a_{1}^{h}, P_{a_{1}^{h}b}, b)$. Then P_{xy}, P_{zw}, P_{ab} are three desired paths in AQ_{n} (see Figure 6(c)).

The lemma holds.



Figure 6. Illustrations of Lemma 3.5.

4. Main results

Theorem 4.1. Let $F \subset V(AQ_n) \cup E(AQ_n)$ with $|F| \le 2n - 3(n \ge 4)$. Then for arbitrary two different vertices u and v in $AQ_n - F$, there exists an error-free Hamiltonian path P_{uv} except (u, v) is a weak vertex-pair in $AQ_n - F$.

Proof. For $|F| \le 2n - 4$, the theorem holds by Lemma 3.1. We only need to consider |F| = 2n - 3.

Now, we prove the theorem by induction on $n \ge 4$. The induction basis for n = 4, 5 holds by computer program (https://github.com/ZhangHeidi/Hypercubes/blob/master/vcn02.c). Supposing that the theorem holds for n - 1 with $n \ge 6$, we must show the theorem holds for n.

We may assume $|F^R| \le |F^L|$ (When $|F^R| \ge |F^L|$, the same proofs apply). Then $|F^R| \le \lfloor \frac{2n-3}{2} \rfloor \le n-1$. Notice that $|N_R(x)| = 2n - 3$ for any vertex $x \in R$, we have $|N_{R-F^R}(x)| \ge n - 2 \ge 4$. Thus $R - F^R$ does not contain weak vertex-pairs.

Case 1. $|F^L| \le 2n - 5$.

Case 1.1. $u, v \in V(L - F^L)$ or $u, v \in V(R - F^R)$. Suppose that $u, v \in V(L - F^L)$.

Case 1.1.1. (u, v) is a *w*-weak vertex-pair in $L - F^L$, *i.e.*, $N_{L-F^L}(w) = \{u, v\}$. Since $|N_L(w)| = 2n - 3$, we have $|F^L| = 2n - 5$ and $|F^R| + |F^C| = 2$. Note that (u, v) is a normal vertex-pair in $AQ_n - F$, we can choose a vertex w^g from $\{w^h, w^c\}$ such that $w^g, ww^g \notin F$.

Because $|V(L) - F^L - \{u, w, v\}| \ge 2^{n-1} - (2n-5) - 3 \ge 22(n \ge 6)$ and $|F^R| + |F^C| = 2$, there is a vertex $y \in V(L) - F^L - \{u, w, v\}$ such that $y^h, yy^h \notin F$ and $y^h \neq w^g$. By Corollary 2.1, (u, y) is a normal vertex-pair in $L - F^L$. By induction hypothesis, a Hamiltonian path P_{uy} exists in $L - F^L$. Notice that $N_{L-F^L}(w) = \{u, v\}$, then $N_{P_{uy}}(w) = \{u, v\}$. Since $|F^R| \le 2 \le 2n - 6$, by Lemma 3.1, a Hamiltonian path $P_{w^gy^h}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^g, P_{w^gy^h}, y^h, y, P_{uy}(y, v), v)$ can therefore be found(see Figure 7(a)).

Case 1.1.2. (u, v) is a normal vertex-pair in $L - F^L$.

Since $|F^L| \le 2n - 5$, by induction hypothesis, a Hamiltonian path P_{uv} exists in $L - F^L$. By $\lfloor \frac{l_{uv}+1}{2} \rfloor - (2n - 3) = \lfloor \frac{2^{n-1}-f_v^L}{2} \rfloor - (2n - 3) \ge 3$, we can choose an edge $ab \in E(P_{uv})$ such that $a^h, b^h, aa^h, bb^h \notin F$.

By induction hypothesis, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uv}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{uv}(b, v), v)$ can therefore be found(see Figure 7(b)).

Case 1.2. $u \in V(L - F^{L})$ and $v \in V(R - F^{R})$.

According to the definition of AQ_n , $|E^C| = 2^n$. Since $|E^C| - 2|F| = 2^n - 2(2n - 3) \ge 46(n \ge 6)$, there is an error-free edge $ab \in E^C$ such that $a, b \notin \{u, v\}, a, b \notin F$ and (u, a) is a normal vertex-pair in $L - F^L$. By induction hypothesis, a Hamiltonian path P_{ua} in $L - F^L$ and a Hamiltonian path P_{bv} in $R - F^R$ exist. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, b, P_{bv}, v)$ can therefore be found(see Figure 7(c)).



Figure 7. Illustrations of Case 1.1 and Case 1.2 of Theorem 4.1.

- **Case 2.** $|F^L| = 2n 4$. Then $|F^R| + |F^C| = 1$.
- **Case 2.1.** $u, v \in V(L F^L)$.

By Lemma 3.3, we can choose an element $f \in F^L$ such that (u, v) is a normal vertex-pair in $L - (F^L - \{f\})$. Let $F_1^L = F^L - \{f\}$. Then $|F_1^L| = 2n - 5$. By induction hypothesis, a Hamiltonian path P_{uv} exists in $L - F_1^L$.

Case 2.1.1. $f \in V(P_{uv}) \cup E(P_{uv})$.

If $f \in E(L) \cap F^{L}$, say f = ab; if $f \in V(L) \cap F^{L}$, say $N_{P_{uv}}(f) = \{a, b\}$; let $P_{uv} = (u, P_{uv}(u, a), a, f, b, P_{uv}(b, v), v)$. Suppose that $|\{a^{h}, b^{h}, aa^{h}, bb^{h}\} \cap F| \leq |\{a^{c}, b^{c}, aa^{c}, bb^{c}\} \cap F|$.

Case 2.1.1.1. $|\{a^h, b^h, aa^h, bb^h\} \cap F| = 0.$

Since $|F^R| \le 1 \le 2n - 6$, by Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uv}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{uv}(b, v), v)$ can therefore be found(see Figure 8(a)).

Case 2.1.1.2. $|\{a^h, b^h, aa^h, bb^h\} \cap F| = 1$. Then $|\{a^c, b^c, aa^c, bb^c\} \cap F| = 1$.

Notice that $|F^R| + |F^C| = 1$, we have $\{a^h, b^h, aa^h, bb^h\} \cap F = \{a^c, b^c, aa^c, bb^c\} \cap F$, *i.e.*, $\{a^h, b^h, aa^h, bb^h\} = \{a^c, b^c, aa^c, bb^c\}$. Then by Proposition 2.1, $a = b^{c_{n-1}}$, $a^h = b^c$ and $a^c = b^h$. Suppose that $a^h \in F$, then $b^h \notin F$. Let $a_1b_1 \in E(P_{uv})$ with $a_1, b_1 \notin \{a, b\}$ and $F_1^R = F^R + \{b^h\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a_1^h b_1^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uv}(u, a), a, b^h, b, P_{uv}(b, a_1), a_1, a_1^h, P_{a_1^h b_1^h}, b_1^h, b_1, P_{uv}(b_1, v), v)$ can therefore be found(see Figure 8(b)).

Case 2.1.2. $f \notin V(P_{uv}) \cup E(P_{uv})$.

Then an edge a_2b_2 can be chosen from P_{uv} with $a_2^h, b_2^h, a_2a_2^h, b_2b_2^h \notin F$. Since $|F^R| \le 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a_2^h b_2^h}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uv}(u, a_2), a_2, a_2^h, P_{a_2^h b_2^h}, b_2^h, b_2, P_{uv}(b_2, v), v)$ can therefore be found(see Figure 8(c)).

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Figure 8. Illustrations of Case 2.1 of Theorem 4.1.

Case 2.2. $u \in V(L - F^{L})$ and $v \in V(R - F^{R})$.

Case 2.2.1. $|F_v^L| \ge 1$. There is a vertex $x \in F_v^L$. Let $F_1^L = F^L - \{x\}$, then $|F_1^L| = |F^L| - 1 = 2n - 5$. **Case 2.2.1.1.** (u, x) is a normal vertex-pair in $L - F_1^L$.

By induction hypothesis, a Hamiltonian path P_{ux} exists in $L - F_1^L$. Let $N_{P_{ux}}(x) = x_1$. Since $|F^R| + |F^C| = 1$, there is a vertex $x_1^g \in \{x_1^h, x_1^c\}$ such that $x_1^g, x_1 x_1^g \notin F$.

(1) $x_1^g = v$. Choose an edge $ab \in E(P_{ux})$ such that $a^h, b^h, aa^h, bb^h \notin F$ and $a, b \notin \{x, x_1\}$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian $P_{uv} = (u, P_{ux}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{ux}(b, x_1), x_1, v)$ can therefore be found(see Figure 9(a)).

(2) $x_1^g \neq v$. Since $|F^R| \leq 1 \leq 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x_1^g v}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux}(u, x_1), x_1, x_1^g, P_{x_1^g v}, v)$ can therefore be found(see Figure 9(b)).



Figure 9. Illustrations of Case 2.2.1.1 of Theorem 4.1.

Case 2.2.1.2. (u, x) is a *w*-weak vertex-pair in $L - F_1^L$. Then $N_{L-F_1^L}(w) = \{u, x\}, i.e., N_{L-F^L}(w) = u$.

By Corollary 2.1., (u, w) is a normal vertex-pair in $L - F_1^L$. Since $|F_1^L| = 2n - 5$, by induction hypothesis, a Hamiltonian path P_{uw} exists in $L - F_1^L$. By $N_{L-F_1^L}(w) = \{u, x\}$, we have $N_{P_{uw}}(w) = \{x\}$. Let $N_{P_{uw}}(x) = \{x_1\}$ with $x_1 \neq w$.

(1) $|\{w^h, w^c\} \cap F| = 0.$

(1.1) $v \in \{w^h, w^c\}$, assume that $v = w^h$. If $x_1 = w^{c_{n-1}}$, then by Proposition 2.1, $x_1^h = w^c, x_1^c = w^h$. We can choose an edge $ab \in E(P_{uw})$ such that $a^h, b^h, aa^h, bb^h \notin F$ and $a, b \notin \{w, x, x_1\}$, then $v, w^c \notin \{a^h, b^h\}$. Let $F_1^R = F^R + \{v, w^c\}$. Then $|F_1^R| \le 3 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{uw}(b, x_1), x_1, w^c, w, v)$ can therefore be found(see Figure 10(a)). If $x_1 \neq w^{c_{n-1}}$, then by Proposition 2.1, $x_1^h \neq w^c, x_1^c \neq w^h$. There exists an error-free vertex $x_1^g \in \{x_1^c, x_1^h\}$ such that $x_1^g, x_1 x_1^g \notin F$. We have $x_1^g \notin \{w^c, v\}$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x_1^g w^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $x_1, x_1, x_1^g, P_{x_1^g w^c}, w^c, w, v$) can therefore be found(see Figure 10(b)).

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Figure 10. Illustrations of Case 2.2.1.2 of Theorem 4.1.

(1.2) $v \notin \{w^h, w^c\}$.

For $x_1 = w^{c_{n-1}}$. By Proposition 2.1, $x_1^h = w^c$, $x_1^c = w^h$, *i.e.*, $v \notin \{x_1^h, x_1^c\}$. Let $F_1^R = F^R + \{w^c\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{w^h v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw}(u, x_1), x_1, w^c, w, w^h, P_{w^h v}, v)$ connecting vertices u and v can therefore be found (see Figure 11(a)).

For $x_1 \neq w^{c_{n-1}}$. By Proposition 2.1, $x_1^h \neq w^c, x_1^c \neq w^h$. There exists a correct vertex $x_1^g \in \{x_1^c, x_1^h\}$ such that $x_1^g, x_1 x_1^g \notin F$. If $x_1^g = v$, we can choose an edge $ab \in E(P_{uw})$ with $a, b \notin \{w, x, x_1\}$, $a^h, aa^h, b^h, bb^h \notin F$ and $a \neq w^{c_{n-1}}, b \neq w^{c_{n-1}}$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \leq 2$. By Lemma 3.4, two disjoint paths $P_{a^hw^h}, P_{w^cb^h}$ with $V(P_{a^hw^h}) \cup V(P_{w^cb^h}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw}(u, a), a, a^h, P_{a^hw^h}, w^h, w, w^c, P_{w^cb^h}, b^h, b, P_{uw}(b, x_1), x_1, v)$ can therefore be found(see Figure 11(b)). If $x_1^g \neq v$. Since $|F^R| \leq 1$, by Lemma 3.4, two disjoint paths $P_{x_1^gw^h}, P_{w^cv}$ with $V(P_{x_1^gw^h}) \cup V(P_{w^cv}) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw}(u, x_1), x_1, x_1^g, P_{x_2^gw^h}, w^h, w, w^c, P_{w^cv}, v)$ can therefore be found(see Figure 11(c)).



Figure 11. Illustrations of Case 2.2.1.2 of Theorem 4.1.

(2) $|\{w^h, w^c\} \cap F| = 1$. Assume that $w^c \in F$. Let $N_{P_{uw}}(u) = \{u_1\}$.

For $x_1 = w^{c_{n-1}}$ or $u_1 = w^{c_{n-1}}$. Assume that $x_1 = w^{c_{n-1}}$, then by Proposition 2.1, $x_1^h = w^c$, $x_1^c = w^h$ *i.e.*, $x_1^h \in F$. Since $|F^c| + |F^R| = 1$, there exists a correct vertex $u_1^g \in \{u_1^c, u_1^h\}$ such that $u_1^g, u_1u_1^g \notin F$ and $u_1^g \neq v$. Let $F_1^R = F^R + \{w^h\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{u_1^g v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^h, x_1, P_{uw}(x_1, u_1), u_1, u_1^g, P_{u_1^g v}, v)$ can therefore be found(see Figure 12(a)).

For $x_1 \neq w^{c_{n-1}}$ and $u_1 \neq w^{c_{n-1}}$. By Proposition 2.1, $x_1^h \neq w^c$, $x_1^c \neq w^h$, $u_1^h \neq w^c$, $u_1^c \neq w^h$. If $v \in \{x_1^h, u_1^h\}$, assume that $v = x_1^h$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \leq 2 \leq 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{w^h u_1^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^h, P_{w^h u_1^h}, u_1^h, u_1, P_{uw}(u_1, x_1), x_1, v)$ can therefore be found(see Figure 12(b)). If $v \notin \{x_1^h, u_1^h\}$, since $|F^R| \leq 1$, by Lemma 3.4, two disjoint paths $P_{w^h x_1^h}, P_{u_1^h v}$ with $V(P_{w^h x_1^h}, \bigcup V(P_{u_1^h v})) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^h, P_{w^h x_1^h}, x_1^h, x_1, P_{uw}(x_1, u_1), u_1, u_1^h, P_{u_1^h v}, v)$ can therefore be found(see Figure 12(c)).

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Figure 12. Illustrations of Case 2.2.1.2 of Theorem 4.1.

Case 2.2.2. $|F_{v}^{L}| = 0$. A faulty edge f can be chosen from F^{L} such that $f \notin E_{L}(v^{h}) \cup E_{L}(v^{c})$. Otherwise, $F^{L} \subset E_{L}(v^{h}) \cup E_{L}(v^{c})$. Consider the following two cases.

Case 2.2.2.1. There is a faulty edge $f \in F^L$ such that $f \notin E(v^h) \cup E(v^c)$. Let $F_1^L = F^L - f$. Then $|F_1^L| = 2n - 5$. By Lemma 3.1, a Hamiltonian cycle *C* containing vertex *u* exists in $L - F_1^L$.

(1) $f \notin E(C)$. Let $a \in N_C(u)$ with $a, aa^h \notin F$ and $C = (u, P_{ua}, a, u)$.

For $a^h \neq v$, since $|F^R| \leq 1$, by Lemma 3.1, a Hamiltonian path $P_{a^h v}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, a^h, P_{a^h v}, v)$ can therefore be found(see Figure 13(a)).

For $a^h = v$, let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. Choose an edge $a_1b_1 \in E(C)$ such that $a_1^h, b_1^h, a_1a_1^h.b_1b_1^h \notin F$ and $a_1, b_1 \notin \{u, a\}$. By Lemma 3.1, a Hamiltonian path $P_{a_1^h b_1^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}(u, a_1), a_1, a_1^h, P_{a_1^h b_1^h}, b_1^h, b_1, P_{ua}(b_1, a), a, v)$ can therefore be found(see Figure 13(b)).

(2) $f \in E(C)$. Let f = ab. Consider the following two cases.

(2.1) $f \in E_C(u)$. Let f = ub and $C = (u, P_{ub}, b, u)$. Since $|F^R| + |F^C| = 1$, there is a correct vertex $b^g \in \{b^h, b^c\}$ such that $bb^g \notin F$. Since $f \notin E(v^h) \cup E(v^c)$, $b^g \neq v$. Then by $|F^R| \leq 1$ and Lemma 3.1, a Hamiltonian path $P_{b^{g_v}}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ub}, b, b^g, P_{b^{g_v}}, v)$ can therefore be found(see Figure 13(c)).



Figure 13. Illustrations of Case 2.2.2.1 of Theorem 4.1.

(2.2) $f \notin E_C(u)$. Let f = ab and $C = (u, P_{ua}, a, b, P_{bu_1}, u_1, u)$. Suppose that $|\{u_1^h, u_1u_1^h, a^h, aa^h, b^h, bb^h\} \cap F| \le |\{u_1^c, u_1u_1^c, a^c, aa^c, b^c, bb^c\} \cap F|$.

 $(2.2.1) |\{u_1^h, u_1u_1^h, a^h, aa^h, b^h, bb^h\} \cap F| = 0.$

For $u_1^h \neq v$, by Lemma 3.4, two disjoint paths $P_{u_1^{h_v}}$, $P_{a^h b^h}$ with $V(P_{u_1^{h_v}}) \cup V(P_{a^h b^h}) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, a^h, P_{a^h b^h}, b^h, b, P_{bu_1}, u_1, u_1^h, P_{u_1^{h_v}}, v)$ can therefore be found(see Figure 14(a)).

For $u_1^h = v$, let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, a^h, P_{a^h b^h}, b^h, b, P_{bu_1}, u_1, v)$ can therefore be found(see Figure 14(b)).

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Figure 14. Illustrations of Case 2.2.2.1 of Theorem 4.1.

 $(2.2.2) |\{u_1^h, u_1u_1^h, a^h, aa^h, b^h, bb^h\} \cap F| = 1. \text{ Then } |\{u_1^c, u_1u_1^c, a^c, aa^c, b^c, bb^c\} \cap F| = 1. \text{ Notice that } |F^R| + |F^C| = 1, \text{ we have } \{u_1^h, u_1u_1^h, a^h, aa^h, b^h, bb^h\} \cap F = \{u_1^c, u_1u_1^c, a^c, aa^c, b^c, bb^c\} \cap F. \text{ Then by Proposition 2.1, } a = b^{c_{n-1}} \text{ or } a = u_1^{c_{n-1}} \text{ or } b = u_1^{c_{n-1}}.$

For $a = b^{c_{n-1}}$, by Proposition 2.1, $a^h = b^c$ and $a^c = b^h$. Suppose that $a^h \in F$. Then $b^h \notin F$ and $|\{u_1^h, u_1u_1^h, u_1^c, u_1u_1^c\} \cap F| = 0$. It follows that a vertex u_1^g can be chosen from $\{u_1^h, u_1^c\}$ such that $u_1^g \neq v$. Let $F_1^R = F^R + \{b^h\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{u_1^g v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, b^h, b, P_{bu_1}, u_1, u_1^g, P_{u_1^g v}, v)$ can therefore be found(see Figure 15(a)).

For $a = u_1^{c_{n-1}}$, by Proposition 2.1, $a^h = u_1^c$ and $a^c = u_1^h$. Suppose that $a^h \in F$. Then $u_1^h \notin F$ and $|\{b^h, bb^h, b^c, bb^c\} \cap F| = 0$. Let $F_1^R = F^R + \{u_1^h\}$. Then $|F_1^R| \le 2 \le 2n - 6$. Notice that $f = ab \notin E(v^h) \cup E(v^c)$, then $v \notin \{b^h, b^c\}$. By Lemma 3.1, a Hamiltonian path $P_{b^h v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua}, a, u_1^h, u_1, P_{u_1b}, b, b^h, P_{b^h v}, v)$ can therefore be found(see Figure 15(b)).

For $b = u_1^{c_{n-1}}$, by Proposition 2.1, $b^h = u_1^c$ and $b^c = u_1^h$. Assume that $u_1^h \in F$. Then $b^h \notin F$ and $|\{a^h, aa^h, a^c, aa^c\} \cap F| = 0$. Let $u_2 = N_C(u)$ with $u_2 \neq u_1$. Then $u_2^h, u_2u_2^h, u_2^c, u_2u_2^c \notin F$. Notice that $f = ab \notin E(v^h) \cup E(v^c)$, then $v \notin \{a^h, a^c\}$. If $u_2 = a^{c_{n-1}}$, since $v \notin \{a^h, a^c\}$, then for any vertex $u_2^g \in \{u_2^h, u_2^c\}$ with $u_2^g \neq v$. If $u_2 \neq a^{c_{n-1}}$, then there is a vertex $u_2^g \in \{u_2^h, u_2^c\}$ such that $u_2^g \neq v$. Since $|F^R| \leq 1$, by Lemma 3.4, two disjoint paths $P_{a^h b^h}, P_{u_2^{g_v}}$ with $V(P_{a^h b^h}) \cup V(P_{u_2^{g_v}}) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ub}, b, b^h, P_{b^h a^h}, a^h, a, P_{ua}(a, u_2), u_2, u_2^g, P_{u_2^{g_v}}, v)$ can therefore be found(see Figure 15(c)). If $u_2 = a$, then we use a^c instead of u_2^g .



Figure 15. Illustrations of Case 2.2.2.1 of Theorem 4.1.

Case 2.2.2. $F^{L} \subset N_{L}(v^{h}) \cup E_{L}(v^{c})$. Suppose that $|F^{L} \cap E_{L}(v^{h})| \ge |F^{L} \cap E_{L}(v^{c})|$. Let $e_{1}, e_{2} \in F^{L} \cap E_{L}(v^{h})$ and $F_{1}^{L} = F^{L} - \{e_{1}, e_{2}\} + \{v^{h}\}$. Then $|F_{1}^{L}| = 2n - 5$. Since $|N_{L}(v^{h})| = 2n - 3$ and $|F^{L}| = 2n - 4$, there is a correct vertex *y* such that $y \in N_{L-F^{L}}(v^{h})$.

(1) $u = v^h$. Choose a vertex x from $V(L) - F^L - \{u, v^c, y\}$ with $x^h, xx^h \notin F$ and $x^h \neq v$. Since $|F^L \cap E_L(v^h)| \ge |F^L \cap E_L(v^c)|$ and $F_1^L = F^L - \{e_1, e_2\} + \{v^h\}$, (x, y) is a normal vertex-pair in $L - F_1^L$. By induction hypothesis, a Hamiltonian path P_{xy} exists in $L - F_1^L$. Since $|F^R| \le 1$, by Lemma 3.1, a Hamiltonian path $P_{x^h v}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, y, P_{yx}, x, x^h, P_{x^h v}, v)$

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can therefore be found(see Figure 16(a)).

(2) $u \neq v^h$.

(2.1) y = u. Then $u \in N_{L-F^{L}}(v^{h})$. Notice that $F^{L} \subset N_{L}(v^{h}) \cup E_{L}(v^{c})$ and $|F^{R}| + |F^{C}| = 1$. Since (u, v) is a normal vertex pair in $AQ_{n} - F$, we have that $v^{h}(v^{h})^{c}$, $(v^{h})^{c} \notin F$. Choose a vertex x from $V(L) - F^{L} - \{v^{h}, v^{c}, y\}$ with $x^{c}, xx^{c} \notin F$ and $x^{c} \neq v$. Since $|F^{L} \cap E_{L}(v^{h})| \geq |F^{L} \cap E_{L}(v^{c})|$ and $F_{1}^{L} = F^{L} - \{e_{1}, e_{2}\} + \{v^{h}\}, (u, x)$ is a normal vertex-pair in $L - F_{1}^{L}$. By induction hypothesis, a Hamiltonian path P_{ux} exists in $L - F_{1}^{L}$.

(2.1.1) $v^h v \notin F$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x^c(v^h)^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux}, x, x^c, P_{x^c(v^h)^c}, (v^h)^c, v^h, v)$ can therefore be found(see Figure 16(b)).



Figure 16. Illustrations of Case 2.2.2.2 of Theorem 4.1.

 $(2.1.2) v^h v \in F$. Let $u_1 \in N_{P_{ux}}(u)$. Since $|F^R| + |F^C| = 1$, $u_1^c, u_1 u_1^c \notin F$.

If $u_1^c = v$, let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n-6$. By Lemma 3.1, a Hamiltonian path $P_{(v^h)^c x^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, v^h, (v^h)^c, P_{(v^h)^c x^c}, x^c, x, P_{ux}(x, u_1), u_1, v)$ connecting vertices u and v can therefore be found(see Figure 17(a)).

If $u_1^c \neq v$, by $|F^R| \leq 1$ and Lemma 3.4, two disjoint paths $P_{(v^h)^c x^c}$, $P_{u_1^c v}$ with $V(P_{(v^h)^c x^c}) \cup V(P_{u_1^c v}) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, v^h, (v^h)^c, P_{(v^h)^c x^c}, x^c, x, P_{ux}(x, u_1), u_1, u_1^c, P_{u_1^c v}, v)$ can therefore be found(see Figure 17(b)).



Figure 17. Illustrations of Case 2.2.2.2 of Theorem 4.1.

(2.2) $y \neq u$. Since $|F^L \cap E_L(v^h)| \geq |F^L \cap E_L(v^c)|$ and $F_1^L = F^L - \{e_1, e_2\} + \{v^h\}$, (u, y) is a normal vertex-pair in $L - F_1^L$. By induction hypothesis, a Hamiltonian path P_{uy} exists in $L - F_1^L$.

If $v^h v \in F$, then $(v^h)^c$, $v^h (v^h)^c \notin F$. Since $|F^R| \le 1$, by Lemma 3.1, a Hamiltonian path $P_{(v^h)^c v}$ exists in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uy}, y, v^h, (v^h)^c)$,

 $P_{(v^h)^c v}, v$) can therefore be found(see Figure 18(a)).

If $v^h v \notin F$, then choose an edge $ab \in E(P_{uy})$ with $a^h, b^h, aa^h, bb^h \notin F$ and $v \notin \{a^h, b^h\}$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uy}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{uy}(b, y), y, v^h, v)$ can therefore be found(see Figure 18(b)).



Figure 18. Illustrations of Case 2.2.2.2 of Theorem 4.1.

Case 2.3. $u, v \in V(R - F^R)$.

Notice that $u^h = (u^c)^{c_{n-1}}$ and $v^h = (v^c)^{c_{n-1}}$. By Proposition 2.2, $|N_L(u^h) \cap N_L(u^c)| = 2$ and $|N_L(v^h) \cap N_L(v^c)| = 2$. Let $N_L(u^h) \cap N_L(u^c) = \{m_1, m_2\}$, $N_L(v^h) \cap N_L(v^c) = \{m_3, m_4\}$, and $F_1^L = \{m_1, m_2, m_3, m_4, u^h u^c, v^h v^c\}$. Since $|F^L| = 2n - 4 \ge 8(n \ge 6)$, $|F^L - F_1^L| \ge 2$. Then we can choose a faulty element f with $f \in F^L - F_1^L$. Then $|F^L - \{f\}| = 2n - 5$. By Lemma 3.1, a Hamiltonian cycle C exists in $L - (F^L - \{f\})$. If C contains f, let $C = (a, P_{ab}, b, f, a)$. If C does not contain f, let $C = (a, P_{ab}, b, a)$ with $ab \notin \{u^h u^c, v^h v^c\}$. Since $|F^R| + |F^C| = 1$, there is a vertex $a^g \in \{a^h, a^c\}$ and a vertex $b^g \in \{b^h, b^c\}$ with $a^g, b^g, aa^g, bb^g \notin F$. If $|\{a^g, b^g\} \cap \{u, v\}| \ge 1$, then by the choice of the above error element $f, a \neq b^{c_{n-1}}$, *i.e.*, a^h, a^c, b^h, b^c are all distinct. Since $|F^R| + |F^C| = 1$, we can choose an error-free set $\{a^g, b^g, aa^g, bb^g\}$ such that $|\{a^g, b^g\} \cap \{u, v\}| = 1$. Therefore, we only need to consider the following two cases.

Case 2.3.1. $|\{a^g, b^g\} \cap \{u, v\}| = 0.$

Case 2.3.1.1. $a^g \neq b^g$.

Since $|F^R| \le 1$, by Lemma 3.4, two disjoint paths P_{ua^g} , $P_{b^{g_v}}$ with $V(P_{ua^g}) \cup V(P_{b^{g_v}}) = V(R - F^R)$ exist in $R - F^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua^g}, a^g, a, P_{ab}, b, b^g, P_{b^{g_v}}, v)$ can therefore be found(see Figure 19(a)).

Case 2.3.1.2. $a^g = b^g$.

Let $F_1^R = F^R + \{a^g\}$. Then $|F_1^R| \le 2$. Choose an edge $a_1b_1 \in E(P_{ab})$ with $a_1, b_1 \notin \{a, b\}$, $a_1^h, b_1^h \notin \{u, v\}$ and $a_1^h, b_1^h, a_1a_1^h, b_1b_1^h \notin F$. By Lemma 3.4, two disjoint paths $P_{ua_1^h}, P_{b_1^h v}$ with $V(P_{ua_1^h}) \cup V(P_{b_1^h v}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ua_1^h}, a_1^h, a_1, P_{ab}(a_1, b), b, a^g, a, P_{ab}(a, b_1), b_1, b_1^h, P_{b_1^h v}, v)$ can therefore be found(see Figure 19(b)).



Figure 19. Illustrations of Case 2.3.1 of Theorem 4.1.

Case 2.3.2. $|\{a^g, b^g\} \cap \{u, v\}| = 1$.

Case 2.3.2.1. $u \in \{a^g, b^g\}$, suppose that $u = a^g$. Let $F_1^R = F^R + \{u\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{b^{g_v}}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, a, P_{ab}, b, b^g, P_{b^{g_v}}, v)$ can therefore be found(see Figure 20(a)). **Case 2.3.2.2.** $v \in \{a^g, b^g\}$, suppose that $v = a^g$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| \le 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path P_{ub^g} exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{ub^g}, b^g, b, P_{ba}, a, v)$ can therefore be found(see Figure 20(b)).



Figure 20. Illustrations of Case 2.3.2 of Theorem 4.1.

Case 3. $|F^L| = 2n - 3$. Then $|F^R| = |F^C| = 0$.

Case 3.1. $u, v \in V(L - F^L)$. By Lemma 3.2, there exist two elements $f_1, f_2 \in F^L$ with (u, v) is a normal vertex-pair in $L - (F^L - \{f_1, f_2\})$. Let $F_1^L = F^L - \{f_1, f_2\}$. Then $|F_1^L| = 2n - 5$. By induction hypothesis, a Hamiltonian path P_{uv} exists in $L - F_1^L$ which may or may not include f_1 or f_2 on it. Removing f_1 and f_2 , path P_{uv} is divided into three, two, or one segments relying on the situations in which f_1 and f_2 are on P_{uv} or not. For the last two situations, we may arbitrarily delete one or two more edges from P_{uv} to make it into three subpaths. Therefore, we may write the path P_{uv} with these subpaths as $(u, P_{uu_1}, u_1, f'_1, x, P_{xy}, y, f'_2, v_1, P_{v_1v}, v)$.

Case 3.1.1. The length of $P_{xy} l_{xy} \ge 1$. By Lemma 3.4, two disjoint paths $P_{u_1^h x^h}$, $P_{y^h v_1^h}$ with $V(P_{u_1^h x^h}) \cup V(P_{y^h v_1^h}) = V(R)$ exist in R. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uu_1}, u_1, u_1^h, P_{u_1^h x^h}, x^h, x, P_{xy}, y, y^h, P_{y^h v_1^h}, v_1^h, v_1, P_{v_1 v}, v)$ can therefore be found(see Figure 21(a)).

Case 3.1.2. The length of $P_{xy} l_{xy} = 0$. Then x = y.

Case 3.1.2.1. $x \neq u_1^{c_{n-1}}$ and $x \neq v_1^{c_{n-1}}$. By Proposition 2.1, $x^h \neq u_1^c$, $x^c \neq u_1^h$ and $x^h \neq v_1^c$, $x^c \neq v_1^h$. By Lemma 3.4, two disjoint paths $P_{u_1^h x^h}$, $P_{x^c v_1^h}$ with $V(P_{u_1^h x^h}) \cup V(P_{x^c v_1^h}) = V(R)$ exist in R. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uu_1}, u_1, u_1^h, P_{u_1^h x^h}, x^h, x, x^c, P_{x^c v_1^h}, v_1^h, v_1, P_{v_1 v}, v)$ can therefore be found(see Figure 21(b)).

Case 3.1.2.2. $x = u_1^{c_{n-1}}$ or $x = v_1^{c_{n-1}}$. Assume that $x = u_1^{c_{n-1}}$. By Proposition 2.1, $x^h = u_1^c$, $x^c = u_1^h$. Let $F_1^R = F^R + \{u_1^h\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x^h v_1^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv}^1 = (u, P_{uu_1}, u_1, u_1^h, x, x^h, P_{x^h v_1^h}, v_1^h, v_1, P_{v_1 v}, v)$ can therefore be found(see Figure 21(c)).



Figure 21. Illustrations of Case 3.1 of Theorem 4.1.

Case 3.2.
$$u \in V(L - F^L), v \in V(R)$$
.

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Case 3.2.1. $|F_v^L| \ge 1$. There exists at least one faulty vertex $x \in F^L$.

Let $F_1^L = F^L - \{x\}$. Then $|F_1^L| = 2n - 4$. By Lemma 3.3, there is an element $f_1 \in F_1^L$ with (u, x) is a normal vertex-pair in $L - (F_1^L - \{f_1\})$. Let $F_2^L = F_1^L - \{f_1\}$. Then $|F_2^L| = 2n - 5$. By induction hypothesis, a Hamiltonian path P_{ux} exists in $L - F_2^L$ which may or may not include f_1 on it. Similar to Case 3.1, we may write the path P_{ux} as $(u, P_{uu_1}, u_1, f'_1, y, P_{yx_1}, x_1, x)$.

Case 3.2.1.1. The length of $P_{yx_1} l_{yx_1} \ge 1$.

(1) $y = u_1^{c_{n-1}}$. By Proposition 2.1, $y^h = u_1^c$ and $y^c = u_1^h$.

(1.1) $v \in \{y^h, u_1^h\}$. We may assume that $u_1^h = v$. We have $v \notin \{y^h, x_1^h\}$. Let $F_1^R = F^R + \{y^h\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x_1^h v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, y^h, y, P_{yx_1}, x_1, x_1^h, P_{x_1^h v}, v)$ can therefore be found(see Figure 22(a)).

(1.2) $v \notin \{y^h, u_1^h\}.$

(1.2.1) $x_1^h \neq v$. By Lemma 3.4, two disjoint paths $P_{u_1^h y^h}$, $P_{x_1^h v}$ with $V(P_{u_1^h y^h}) \cup V(P_{x_1^h v}) = V(R)$ exist in *R*. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, u_1^h, P_{u_1^h y^h}, y^h, y, P_{yx_1}, x_1, x_1^h, P_{x_1^h v}, v)$ can therefore be found(see Figure 22(b)).

(1.2.2) $x_1^h = v$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{u_1^h y^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, u_1^h, P_{u_1^h y^h}, y^h, y, P_{yx_1}, x_1, v)$ can therefore be found(see Figure 22(c)).



Figure 22. Illustrations of Case 3.2.1.1 of Theorem 4.1.

(2) $y \neq u_1^{c_{n-1}}$. By Proposition 2.1, $y^h \neq u_1^c$ and $y^c \neq u_1^h$. There exists $\{u_1^g, y^g\} \in \{\{u_1^h, y^h\}, \{u_1^c, y^c\}\}$ such that $v \notin \{u_1^g, y^g\}$.

For $x_1^h \neq v$, the proof is similar to (1.2.1) of Case 3.2.1.1.

For $x_1^h = v$, the proof is similar to (1.2.2) of Case 3.2.1.1.

Case 3.2.1.2. The length of $P_{yx_1} l_{yx_1} = 0$. Then $x_1 = y$.

(1) $y = u_1^{c_{n-1}}$. By Proposition 2.1, $y^h = u_1^c$ and $y^c = u_1^h$.

(1.1) $v \in \{y^h, u_1^h\}$. We may assume that $v = y^h$. Choose an edge $ab \in P_{ux}$ with $a, b \notin \{u_1, y, x\}$. Then $v \notin \{a^h, b^h\}$. Let $F_1^R = F^R + \{v, u_1^h\}$. Then $|F_1^R| = 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}(u, a), a, a^h, P_{a^h b^h}, b^h, b, P_{uu_1}(b, u_1), u_1, u_1^h, y, v)$ can therefore be found(see Figure 23(a)).

(1.2) $v \notin \{y^h, u_1^h\}$. Let $F_1^R = F^R + \{u_1^h\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{y^h v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, u_1^h, y, y^h, P_{y^h v}, v)$ can therefore be found(see Figure 23(b)).

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Figure 23. Illustrations of Case 3.2.1.2 of Theorem 4.1.

(2) $y \neq u_1^{c_{n-1}}$. By Proposition 2.1, $y^h \neq u_1^c$ and $y^c \neq u_1^h$. There is a vertex $u_1^g \in \{u_1^h, u_1^c\}$ such that $u_1^g \neq v$.

(2.1) $v \in \{y^h, y^c\}$. We may assume that $v = y^h$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{u_1^g y^c}$ exists in $R - F_1^R$. Let $P_{uv} = (u, P_{uu_1}, u_1, u_1^g, P_{u_1^g y^c}, y^c, y, v)$. An error-free Hamiltonian path P_{uv} can therefore be found(see Figure 24(a)).

(2.2) $v \notin \{y^h, y^c\}$. By Lemma 3.4, two disjoint paths $P_{u_1^g y^c}$, $P_{y^h v}$ with $V(P_{u_1^g y^c}) \cup V(P_{y^h v}) = V(R)$ exist in *R*. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, u_1^g, P_{u_1^g y^c}, y^c, y, y^h, P_{v^h v}, v)$ can therefore be found(see Figure 24(b)).



Figure 24. Illustrations of Case 3.2.1.2 of Theorem 4.1.

Case 3.2.2. $|F_v^L| = 0.$

Case 3.2.2.1. For any vertex $y \in V(L)$, $|E_L(y) \cap F^L| \le 2$.

Since $|F^L| = 2n - 3 \ge 9$ ($n \ge 6$), we can choose two edges $f_1, f_2 \in F^L$ with $f_1, f_2 \notin E_L(v^h) \cup E_L(v^c)$ and f_1, f_2 do not share the same vertex. Let $F_1^L = F^L - \{f_1, f_2\}$. Then $|F_1^L| = 2n - 5$. We may assume that $v^h \ne u$. By induction hypothesis, there exists a Hamiltonian path P_{uv^h} in $L - F_1^L$ which may or may not include f_1 or f_2 on it. Removing f_1 and f_2 , path P_{uv^h} is divided into three, two, or one segments relying on the situations in which f_1 and f_2 are on P_{uv^h} or not. For the last two situations, we may delete one or two more edges which are not incident to vertices v^h, v^c from P_{uv^h} to make it into three subpaths. Therefore, we may write the path P_{uv^h} with these subpaths as $(u, P_{uu_1}, u_1, f'_1, x, P_{xy}, y, f'_2, v_1, P_{v_1v^h}, v^h)$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths $P_{u_1^h x^h}$ and $P_{y^h v_1^h}$ with $V(P_{u_1^h x^h}) \cup V(P_{y^h v_1^h}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uu_1}, u_1, u_1^h, P_{u_1^h x^h}, x^h, x, P_{xy}, y, y^h, P_{y^h v_1^h}, v_1^h, v_1, P_{v_1v^h}, v^h, v)$ can therefore be found(see Figure 25(a)).

Case 3.2.2. A vertex *y* exists in *L* with $|E_L(y) \cap F^L| \ge 3$.

Let $e_1, e_2, e_3 \in E_L(y) \cap F^L$ and $F_1^L = F^L - \{e_1, e_2, e_3\} + y$. Then $|F_1^L| = 2n - 5$.

(1) y = u. There is a vertex $u^g \in \{u^h, u^c\}$ with $u^g \neq v$. Let $N_L(u^g) = \{u, u_1\}$ and $z \in V(L) - \{u, u_1, v^h, v^c\}$ with (u_1, z) is a normal vertex-pair in $L - F_1^L$. Then $z^h \neq v$. By induction hypothesis, a Hamiltonian path P_{u_1z} exists in $L - F_1^L$. Let $F_1^R = F^R + \{u^g\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path

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 $P_{z^{h_v}}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, u^g, u_1, P_{u_1z}, z, z^h, P_{z^{h_v}}, v)$ can therefore be found(see Figure 25(b)).



Figure 25. Illustrations of Case 3.2.2.1 and Case 3.2.2.2 of Theorem 4.1.

(2) $y \neq u$.

(2.1) $y \in \{v^h, v^c\}$. We may assume that $y = v^c$. Let $z \in V(L) - \{u, v^h, v^c\}$ with (u, z) is a normal vertex-pair in $L - F_1^L$. Then $z^h \neq v$. By induction hypothesis, a Hamiltonian path P_{uz} exists in $L - F_1^L$. Let $F_1^R = F^R + \{v\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{z^h(v^c)^h}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uz}, z, z^h, P_{z^h(v^c)^h}, (v^c)^h, v^c, v)$ can therefore be found(see Figure 26(a)).

(2.2) $y \notin \{v^h, v^c\}$. Let $z \in V(L) - \{u, y, v^h, v^c\}$ with (u, z) is a normal vertex-pair in $L - F_1^L$ and $z \neq y^{c_{n-1}}$. By Proposition 2.1, $z^h \neq y^c$ and $z^c \neq y^h$. By induction hypothesis, a Hamiltonian path P_{uz} exists in $L - F_1^L$. By Lemma 3.4, two disjoint paths $P_{z^hy^h}$ and P_{y^cv} with $V(P_{z^hy^h}) \cup V(P_{y^cv}) = V(R)$ exist in R. An error-free Hamiltonian path $P_{uv} = (u, P_{uz}, z, z^h, P_{z^hy^h}, y^h, y, y^c, P_{y^cv}, v)$ can therefore be found(see Figure 26(b)).



Figure 26. Illustrations of Case 3.2.2.2 of Theorem 4.1.

Case 3.3. $u, v \in V(R)$.

Case 3.3.1. $|F_{\nu}^{L}| \ge 2$. There are at least two vertices $x_1, x_2 \in F^{L}$. Let $F_1^{L} = F^{L} - \{x_1, x_2\}$. Then $|F_1^{L}| = 2n - 5$.

Case 3.3.1.1. (x_1, x_2) is a normal vertex-pair in $L - F_1^L$. By induction hypothesis, a Hamiltonian path $P_{x_1x_2}$ exists in $L - F_1^L$. Let $N_{P_{x_1x_2}}(x_1) = x_{11}$ and $N_{P_{x_1x_2}}(x_2) = x_{21}$.

(1) There is a vertex $x_{11}^g \in \{x_{11}^h, x_{11}^c\}$ and a vertex $x_{21}^g \in \{x_{21}^h, x_{21}^c\}$ such that $x_{11}^g, x_{21}^g \notin \{u, v\}$. By Lemma 3.4, two disjoint paths $P_{ux_{11}^g}, P_{x_{21}^g v}$ with $V(P_{ux_{11}^g}) \cup V(P_{x_{21}^g v}) = V(R)$ exist in R. An error-free Hamiltonian path $P_{uv} = (u, P_{ux_{11}^g}, x_{11}^g, x_{11})$,

 $P_{x_1x_2}(x_{11}, x_{21}), x_{21}, x_{21}^g, P_{x_{21}^g, v}, v)$ can therefore be found(see Figure 27(a)).

(2) For any vertex $x_{11}^{g} \in \{x_{11}^h, x_{11}^c\}$ and vertex $x_{21}^g \in \{x_{21}^h, x_{21}^c\}, \{x_{11}^g, x_{21}^g\} = \{u, v\}$. We assume that $x_{11}^g = u$ and $x_{21}^g = v$. Let $ab \in E(P_{x_1x_2}(x_{11}, x_{21}))$ with $a, b \notin \{x_{11}, x_{21}\}$ and $F_1^R = F^R + \{u, v\}$. Then $|F_1^R| = 2 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{a^h b^h}$ exists in $R - F_1^R$. An error-free Hamiltonian

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path $P_{uv} = (u, x_{11}, P_{x_1x_2}(x_{11}, a), a, a^h, P_{a^hb^h}, b^h, b, P_{x_1x_2}(b, x_{21}), x_{21}, v)$ can therefore be found(see Figure 27(b)).

(3) $\{x_{11}^h, x_{11}^c\} = \{u, v\}$ and $u, v \notin \{x_{21}^h, x_{21}^c\}$ or $\{x_{21}^h, x_{21}^c\} = \{u, v\}$ and $u, v \notin \{x_{11}^h, x_{11}^c\}$. Assume that $\{x_{11}^h, x_{11}^c\} = \{u, v\}$ and $u, v \notin \{x_{21}^h, x_{21}^c\}$. Let $F_1^R = F^R + \{u\}$. Then $|F_1^R| = 1 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{x_{21}^h v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, x_{11}, P_{x_{1x_2}}(x_{11}, x_{21}), x_{21}, x_{21}^h, P_{x_{11}^h v}, v)$ can therefore be found(see Figure 27(c)).



Figure 27. Illustrations of Case 3.3.1.1 of Theorem 4.1.

Case 3.3.1.2. (x_1, x_2) is a *w*-weak vertex-pair in $L - F_1^L$, *i.e.*, $F^L \subset N_L(w) \cup E_L(w)$, $|N_{L-F^L}(w)| = 0$ and $N_{L-F_1^L}(w) = \{x_1, x_2\}$. Since (u, v) is a normal vertex-pair in $AQ_n - F$, $|\{w^h, w^c\} \cap \{u, v\}| \le 1$. By Corollary 2.1, (w, x_2) is a normal vertex-pair in $L - F_1^L$. By induction hypothesis, a Hamiltonian path P_{wx_2} exists in $L - F_1^L$. Since $N_{L-F_1^L}(w) = \{x_1, x_2\}$, we have $N_{P_{wx_2}}(w) = x_1$.

(1) $|\{w^h, w^c\} \cap \{u, v\}| = 1$. Suppose that $u \in \{w^h, w^c\}$, say $u = w^h$. Let $N_{P_{wx_2}}(x_1) - w = x_{11}$ and $N_{P_{wx_2}}(x_2) = x_{21}$.

(1.1) $x_{11} = w^{c_{n-1}}$ or $x_{21} = w^{c_{n-1}}$. Suppose that $x_{11} = w^{c_{n-1}}$. By Proposition 2.1, $w^h = x_{11}^c$ and $w^c = x_{11}^h$. There is a vertex $x_{21}^g \in \{x_{21}^h, x_{21}^c\}$ such that $x_{21}^g \neq v$. Let $F_1^R = F^R + \{u, w^c\}$. Then $|F_1^R| = 2 \leq 2n - 6$. By Lemma 3.1, an error-free Hamiltonian path $P_{x_{21}^g v}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^c, x_{11}, P_{wx_2}(x_{11}, x_{21}), x_{21}, x_{21}^g, P_{x_{21}^g v}, v)$ can therefore be found(see Figure 28(a)).

(1.2) $x_{11} \neq w^{c_{n-1}}$ and $x_{21} \neq w^{c_{n-1}}$. Then by Proposition 2.1, $x_{11}^h \neq w^c$, $x_{11}^c \neq w^h$ and $x_{21}^c \neq w^h$, $x_{21}^h \neq w^c$.

If $v \in \{x_{21}^c, x_{11}^c\}$, we may assume that $x_{21}^c = v$. Let $F_1^R = F^R + \{u, v\}$. Then $|F_1^R| = 2 \le 2n - 6$. By Lemma 3.1, an error-free Hamiltonian path $P_{w^c x_{11}^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^c, P_{w^c x_{11}^c}, x_{11}^c, x_{11}, P_{wx_2}(x_{11}, x_{21}), x_{21}, v)$ can therefore be found(see Figure 28(b)).

If $v \notin \{x_{21}^c, x_{11}^c\}$, let $F_1^R = F^R + \{u\}$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths $P_{w^c x_{11}^c}$, $P_{x_{21}^c v}$ with $V(P_{w^c x_{11}^c}) \cup V(P_{x_{21}^c v}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, w, w^c, P_{w^c x_{11}^c}, x_{11}^c, x_{11}, P_{wx_2}(x_{11}, x_{21}), x_{21}, x_{21}^c, P_{x_{21}^c v}, v)$ can therefore be found(see Figure 28(c)).



Figure 28. Illustrations of Case 3.3.1.2 of Theorem 4.1.

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(2) $|\{w^h, w^c\} \cap \{u, v\}| = 0.$

(2.1) $x_{11} = w^{c_{n-1}}$ or $x_{21} = w^{c_{n-1}}$. Assume that $x_{11} = w^{c_{n-1}}$. By Proposition 2.1, $w^h = x_{11}^c$ and $w^c = x_{11}^h$. If at least one vertex of $\{u, v\}$, say v, is adjacent to x_{21} . Let $F_1^R = F^R + \{w^c, v\}$. Then $|F_1^R| = 2 \le 2n - 6$. By Lemma 3.1, an error-free Hamiltonian path P_{uw^h} exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw^h}, w^h, w, w^c, x_{11}, P_{wx_2}(x_{11}, x_{21}), x_{21},$

v) can therefore be found(see Figure 29(a)).

If both of *u* and *v* are not adjacent to x_{21} . Let $F_1^R = F^R + w^c$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths $P_{uw^h}, P_{x_{21}^h v}$ with $V(P_{uw^h}) \cup V(P_{x_{21}^h v}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw^h}, w^h, w, w^c, x_{11}, P_{wx_2}(x_{11}, x_{21}),$

 $x_{21}, x_{21}^h, P_{x_{21}^h, \nu}, \nu$) can therefore be found(see Figure 29(b)).



Figure 29. Illustrations of Case 3.3.1.2 of Theorem 4.1.

(2.2) $x_{11} \neq w^{c_{n-1}}$ and $x_{21} \neq w^{c_{n-1}}$. Then by Proposition 2.1, $x_{11}^h \neq w^c$, $x_{11}^c \neq w^h$ and $x_{11}^c \neq w^h$, $x_{11}^h \neq w^c$.

(2.2.1) There is a vertex $x_{11}^g \in \{x_{11}^h, x_{11}^c\}$ with $x_{11}^g \notin \{u, v\}$ and $x_{21}^g \in \{x_{21}^h, x_{21}^c\}$ with $x_{21}^g \notin \{u, v\}$. By Lemma 3.5, three disjoint paths $P_{uw^h}, P_{w^c x_{11}^g}, P_{x_{21}^g v}$ exist in R with $V(P_{uw^h}) \cup V(P_{w^c x_{11}^g}) \cup V(P_{x_{21}^g v}) = V(R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{uw^h}, w^h, w, w^c, P_{w^c x_{11}^g}, x_{11}^g, x_{11}, P_{wx_2}(x_{11}, x_{21}), x_{21}, x_{21}^g, P_{x_{21}^g v}, v)$ can therefore be found(see Figure 30(a)).

 $(2.2.2) \{x_{11}^{h}, x_{11}^{c}\} = \{u, v\} \text{ and } u, v \notin \{x_{21}^{h}, x_{21}^{c}\} \text{ or } \{x_{21}^{h}, x_{21}^{c}\} = \{u, v\} \text{ and } u, v \notin \{x_{11}^{h}, x_{11}^{c}\}.$ Assume that $\{x_{11}^{h}, x_{11}^{c}\} = \{u, v\}$ and $u, v \notin \{x_{21}^{h}, x_{21}^{c}\}.$ Let $F_{1}^{R} = F^{R} + \{u\}.$ Then $|F_{1}^{R}| = 1$. By Lemma 3.4, two disjoint paths $P_{x_{21}^{h}w^{h}}, P_{w^{c}v}$ exist in $R - F_{1}^{R}$ with $V(P_{x_{21}^{h}w^{h}}) \cup V(P_{w^{c}v}) = V(R - F_{1}^{R}).$ An error-free Hamiltonian path $P_{uv} = (u, x_{11}, P_{wx_{2}}(x_{11}, x_{21}), x_{21}, x_{21}^{h}, P_{x_{21}^{h}w^{h}}, w^{h}, w, w^{c}, P_{w^{c}v}, v)$ can therefore be found(see Figure 30(b)).



Figure 30. Illustrations of Case 3.3.1.2 of Theorem 4.1.

(2.2.3) $\{x_{21}^h, x_{21}^c, x_{11}^h, x_{11}^c\} = \{u, v\}$. We may assume that $x_{21}^h = u$ and $x_{11}^h = v$. Let $a \in N_{P_{wx_2}}(x_{11})$ with $a \neq x_1$ and $b \in N_{P_{wx_2}}(a)$ with $b \neq x_{11}$.

(2.2.3.1) $a = w^{c_{n-1}}$. Then $a^h = w^c$ and $a^c = w^h$. Let $F_1^R = F^R + \{u, v, w^h\}$. Then $|F_1^R| = 3 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{b^c w^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} =$

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 $(u, x_{21}, P_{wx_2}(x_{21}, b), b, b^c, P_{b^c w^c}, w^c, w, w^h, a, x_{11}, v)$ can therefore be found(see Figure 31(a)).

(2.2.3.2) $b = w^{c_{n-1}}$. Then $b^h = w^c$ and $b^c = w^h$. Let $F_1^R = F^R + \{u, v, w^h\}$. Then $|F_1^R| = 3 \le 2n - 6$. By Lemma 3.1, a Hamiltonian path $P_{w^c a^c}$ exists in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, x_{21}, P_{wx_2}(x_{21}, b), b, w^h, w, w^c, P_{w^c a^c}, a^c, a, x_{11}, v)$ can therefore be found(see Figure 31(b)).

(2.2.3.3) $a \neq w^{c_{n-1}}$ and $b \neq w^{c_{n-1}}$. Then $a^h \neq w^c$, $a^c \neq w^h$ and $b^h \neq w^c$, $b^c = w^h$. Let $F_1^R = F^R + \{u, v\}$. Then $|F_1^R| = 2$. By Lemma 3.4, two disjoint paths $P_{b^h w^h}$, $P_{w^c a^h}$ exist in $R - F_1^R$ with $V(P_{b^h w^h}) \cup V(P_{w^c a^h}) = V(R - F_1^R)$. An error-free Hamiltonian path $P_{uv} = (u, x_{21}, P_{wx_2}(x_{21}, b), b, b^h, P_{b^h w^h}, w^h, w, w^c, P_{w^c a^h}, a^h, a, x_{11}, v)$ can therefore be found(see Figure 31(c)).



Figure 31. Illustrations of Case 3.3.1.2 of Theorem 4.1.

Case 3.3.2. $|F_v^L| \le 1$.

Case 3.3.2.1. For any vertex x with $|E_L(x) \cap F^L| \le 2$. Since $|F^L \cap E(L)| \ge 2n - 4 \ge 8(n \ge 6)$, we can choose two edges $e_1, e_2 \notin E_L(u^c) \cup E_L(v^h) \cup E_L(v^c)$. Let $F_1^L = F^L - \{e_1, e_2\}$. Then $|F_1^L| = 2n - 5$. By Lemma 3.1, a Hamiltonian cycle C will exist in $L - F_1^L$ which may or may not include e_1 or e_2 on it. Removing e_1 and e_2 , cycle C is divided into two, one or zero pieces depending on the cases in which e_1 and e_2 are on C or not. For the last two cases, we may choose to delete one or two more edges which are not incident to vertices u^h, u^c, v^h, v^c from C to make it into two subpaths. Therefore, we may write cycle C using these subpaths as $(x, e_1', x_1, P_{x_1y_1}, y_1, e_2', y, P_{yx}, x)$.

(1) The length of $P_{yx} l_{yx} \ge 1$. Since $|F^R| = 0$, by Lemma 3.5, three disjoint paths $P_{ux_1^c}$, $P_{y_1^cy^c}$, $P_{x^{c_v}}$ exist in R with $V(P_{ux_1^c}) \cup V(P_{y_1^cy^c}) \cup V(P_{x^{c_v}}) = V(R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux_1^c}, x_1^c, x_1, P_{x_1y_1}, y_1, y_1^c, P_{y_1^cy^c}, y^c, y, P_{yx}, x, x^c, P_{x^{c_v}}, v)$ can therefore be found(see Figure 32(a)).

(2) The length of $P_{yx} l_{yx} = 0$. Then x = y.

(2.1) $x = u^h$. Let $F_1^R = F^R + \{u\}$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths $P_{x^c y_1^c}$, $P_{x_1^c y}$ with $V(P_{x^c y_1^c}) \cup V(P_{x_1^c y}) = V(R - F_1^R)$ exist in $R - F_1^R$. An error-free Hamiltonian path $P_{uv} = (u, x, x^c, P_{x^c y_1^c}, y_1^c, y_1, P_{y_1 x_1}, x_1, x_1^c, P_{x_1^c v}, v)$ can therefore be found(see Figure 32(b)).



Figure 32. Illustrations of Case 3.3.2.1 of Theorem 4.1.

(2.2) $x \neq u^{h}$.

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(2.2.1) $x \neq x_1^{c_{n-1}}$ and $x \neq y_1^{c_{n-1}}$. By Proposition 2.1, $x^h \neq x_1^c$, $x^c \neq x_1^h$ and $x^h \neq y_1^c$, $x^c \neq y_1^h$. By Lemma 3.5, three disjoint paths $P_{ux_1^c}$, $P_{y_1^cx^c}$ and P_{x^hv} exist in R with $V(P_{ux_1^c}) \cup V(P_{y_1^cx^c}) \cup V(P_{x^hv}) = V(R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux_1^c}, x_1^c, x_1, P_{x_1y_1}, y_1, y_1^c, P_{y_1^cx^c}, x^c, x, x^h, P_{x^hv}, v)$ can therefore be found(see Figure 33(a)).

(2.2.2) $x = x_1^{c_{n-1}}$ or $x = y_1^{c_{n-1}}$. Suppose that $x = x_1^{c_{n-1}}$. By Proposition 2.1, $x^h = x_1^c$ and $x^c = x_1^h$. Let $F_1^R = F^R + \{x_1^c\}$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths P_{ux^c} , $P_{y_1^cv}$ exist in $R - F_1^R$ with $V(P_{ux^c}) \cup V(P_{y_1^cv}) = V(R - F_1^R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux^c}, x^c, x, x_1^c, x_1, P_{x_1y_1}, y_1, y_1^c, P_{y_1^cv}, v)$ can therefore be found(see Figure 33(b)).



Figure 33. Illustrations of Case 3.3.2.1 of Theorem 4.1.

Case 3.3.2.2. There exists a vertex x with $|E_L(x) \cap F^L| \ge 3$. Let $e_1, e_2, e_3 \in E_L(x) \cap F^L$ and $F_1^L = F^L - \{e_1, e_2, e_3\} + x$. Then $|F_1^L| = 2n - 5$.

(1) There exists no correct element incident to vertex x in $L - F^L$. Since (u, v) is a normal vertex-pair in $AQ_n - F$, $|N_{AQ_n - F}(x) \cap \{u, v\}| \le 1$. Choose two different vertices x_1, y from $V(L) - \{x, u^h, u^c, v^h, v^c\}$ with $x_1 \ne x^{c_{n-1}}, y \ne x^{c_{n-1}}$ and (x_1, y) is a normal vertex-pair in $L - F_1^L$. By induction hypothesis, a Hamiltonian path P_{x_1y} exists in $L - F_1^L$. Notice that $x_1, y \in V(L) - \{x, u^h, u^c, v^h, v^c\}$, then $u, v \notin \{x_1^h, x_1^c, y^h, y^c\}$.

(1.1) $|N_{AQ_n-F}(x) \cap \{u,v\}| = 1$. Assume that $u \in N_{AQ_n-F}(x)$ and $u = x^h$. Let $F_1^R = F^R + \{u\}$. Then $|F_1^R| = 1$. By Lemma 3.4, two disjoint paths $P_{x^cx_1^c}$, P_{y^cv} exist in $R - F_1^R$ with $V(P_{x^cx_1^c}) \cup V(P_{y^cv}) = V(R - F_1^R)$. An error-free Hamiltonian path $P_{uv} = (u, x, x^c, P_{x^cx_1^c}, x_1^c, x_1, P_{x_1y}, y, y^c, P_{y^cv}, v)$ can therefore be found(see Figure 34(a)).

(1.2) $|N_{AQ_n-F}(x) \cap \{u,v\}| = 0$. By Lemma 3.5, three disjoint paths P_{ux^h} , $P_{x^cx_1^c}$, P_{y^cv} exist in R with $V(P_{ux^h}) \cup V(P_{x^cx_1^c}) \cup V(P_{y^cv}) = V(R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux^h}, x^h, x, x^c, P_{x^cx_1^c}, x_1^c, x_1, P_{x_1y}, y, y^c, P_{y^cv}, v)$ can therefore be found(see Figure 34(b)).

(2) There is a vertex x_1 incident to vertex x in $L - F^L$. Choose a vertex $y \in V(L) - \{x, x_1, v^h, v^c, u^h, u^c\}$ with (x_1, y) is a normal vertex-pair in $L - F_1^L$, then $u, v \notin \{y^h, y^c\}$. By induction hypothesis, a Hamiltonian path P_{x_1y} exists in $L - F_1^L$. There is a vertex $x^g \in \{x^h, x^c\}$ with $x^g \neq \{u, v\}$. By Lemma 3.4, two disjoint paths P_{ux^g} , P_{y^cv} exist in R with $V(P_{ux^g}) \cup V(P_{y^cv}) = V(R)$. An error-free Hamiltonian path $P_{uv} = (u, P_{ux^g}, x^g, x, x_1, P_{x_1y}, y, y^c, P_{y^cv}, v)$ can therefore be found(see Figure 34(c)).



Figure 34. Illustrations of Case 3.3.2.2 of Theorem 4.1.

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Combining the above cases, we completed the proof of the Theorem 4.1.

5. Conclusions

This paper studied the Hamiltonian path in *n*-dimensional augmented cube AQ_n with a set *F* of up to 2n - 3 faulty elements. We have proved that for arbitrary vertex-pair (u, v) in $AQ_n - F$, there exists a fault-free Hamiltonian path that joins vertices *u* and *v* with the exception of (u, v), which is a weak vertex-pair in $AQ_n - F(n \ge 4)$. It is worth pointing out that we also proved that if there is a weak vertex-pair in $AQ_n - F$, there is at most one pair. This paper improved the current result that AQ_n is 2n - 4 fault-tolerant Hamiltonian connected. Since the degree of each vertex is 2n - 1 in AQ_n , our result is optimal and sharp under the condition of no restriction to each vertex.

The result of the paper can be further improved. One possible research is the Hamiltonian connectivity when the correct degree of each vertex is restricted and the fault-tolerant bound of AQ_n may be improved; another is the issue of finding the shorter path under the current optimal fault-tolerant bound in which when $|F| \le 2n - 3$, any correct two vertices u, v in $AQ_n - F$ have a path of each length from d-rank to $2^n - f_v - 2$ connecting them, where d is the distance of u, v and f_v is the number of faulty vertices in AQ_n .

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Conflict of interest

The authors declare that they have no competing interests.

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