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Research article

On the metric basis in wheels with consecutive missing spokes

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Abstract: If *G* is a connected graph, the *distance* d(u, v) between two vertices $u, v \in V(G)$ is the length of a shortest path between them. Let $W = \{w_1, w_2, \ldots, w_k\}$ be an ordered set of vertices of *G* and let *v* be a vertex of *G*. The *representation* r(v|W) of *v* with respect to *W* is the k-tuple $(d(v, w_1), d(v, w_2), \ldots, d(v, w_k))$. *W* is called a *resolving set* or a *locating set* if every vertex of *G* have distinct representations with respect to *W*. A resolving set of minimum cardinality is called a *metric basis* for *G* and this cardinality is the *metric dimension* of *G*, denoted by $\beta(G)$. The metric dimension of some wheel related graphs is studied recently by Siddiqui and Imran. In this paper, we study the metric dimension of wheels with *k* consecutive missing spokes denoted by W(n, k). We compute the exact value of the metric dimensions. It is natural to ask for the characterization of graphs with an unbounded metric dimension. The exchange property of the bases in a vector space does not hold for minimal resolving sets of wheels with *k*-consecutive missing spokes denoted by W(n, k).

Keywords: metric dimension; basis; resolving set; wheel; missing spokes; exchange property **Mathematics Subject Classification:** 05C12

1. Introduction and preliminary results

Metric dimension is an important parameter in metric graph theory that has appeared in numerous applications of graph theory, as diverse as, facility location problems, pharmaceutical chemistry [5,6], long range aids to navigation, navigation of robots in networks [17], combinatorial optimization [21]

and sonar and coast guard Loran [23], to name a few. Metric dimension is a generalization of affine dimension to arbitrary metric spaces (provided a resolving set exists).

In a connected graph G, the distance d(u, v) between two vertices $u, v \in V(G)$ is the length of a shortest path between them. Let $W = \{w_1, w_2, \ldots, w_k\}$ be an ordered set of vertices of G and let v be a vertex of G. The representation r(v|W) of v with respect to W is the k-tuple $(d(v, w_1), d(v, w_2), d(v, w_3), \ldots, d(v, w_k))$. W is called a resolving set [6] or locating set [23] if every vertex of G is uniquely identified by its distances from the vertices of W, or equivalently, if distinct vertices of G have distinct representations with respect to W. A resolving set of minimum cardinality is called a basis for G and this cardinality is the metric dimension of G, denoted by $\beta(G)$ [3].

For a given ordered set of vertices $W = \{w_1, w_2, ..., w_k\}$ of a graph *G*, the *i*th component of r(v|W) is 0 if and only if $v = w_i$. Thus, to show that *W* is a resolving set it is sufficient to verify that $r(x|W) \neq r(y|W)$ for each pair of distinct vertices $x, y \in V(G) \setminus W$.

A useful property in computing the metric dimension denoted by $\beta(G)$ is the following lemma:

Lemma 1.1. [25] Let W be a resolving set for a connected graph G and $u, v \in V(G)$. If d(u, w) = d(v, w) for all vertices $w \in V(G) \setminus \{u, v\}$, then $\{u, v\} \cap W \neq \emptyset$.

Motivated by the problem of uniquely determining the location of an intruder in a network, the concept of metric dimension was first introduced by Slater in [23, 24] and studied independently by Harary and Melter in [9]. Applications of this invariant to the navigation of robots in networks are discussed in [17] and applications to chemistry in [6] while applications to problem of pattern recognition and image processing, some of which involve the use of hierarchical data structures are given in [18].

Let \mathcal{F} be a family of connected graphs $G_n : \mathcal{F} = (G_n)_{n\geq 1}$ depending on *n* as follows: the order $|V(G)| = \varphi(n)$ and $\lim_{n\to\infty} \varphi(n) = \infty$. If there exists a constant C > 0 such that $\beta(G_n) \leq C$ for every $n \geq 1$ then we shall say that \mathcal{F} has bounded metric dimension; otherwise \mathcal{F} has an unbounded metric dimension.

If all graphs in \mathcal{F} have the same metric dimension (which does not depend on *n*), \mathcal{F} is called a family with constant metric dimension [15]. A connected graph *G* has $\beta(G) = 1$ if and only if *G* is a *path* [6]; *cycles* C_n have metric dimension 2 for every $n \ge 3$. Also generalized Petersen graphs P(n, 2), *antiprisms* A_n and *circulant graphs* $C_n(1, 2)$ are families of graphs with constant metric dimension have been studied in [13, 20]. An example of a family which has bounded metric dimension is the family of *prisms*. Also generalized Petersen graphs P(n, 3) have bounded metric dimension [10].

Note that the problem of determining whether $\beta(G) < k$ is an *NP*-complete problem [8]. Some bounds for this invariant, in terms of the diameter of the graph, are given in [17] and it was shown in [6,17–19] that the metric dimension of *trees* can be determined efficiently. It appears unlikely that significant progress can be made in determining the dimension of a graph unless it belongs to a class for which the distances between vertices can be described in some systematic manner.

The metric dimension of families of graph having unbounded metric dimension is an interesting problem in the theory of resolving set. In [3], it was shown that the wheel graph W_n have unbounded metric dimension with $\beta(W_n) = \lfloor \frac{2n+2}{5} \rfloor$ for every $n \ge 7$. Later, Javaid et al. [26] show that the graph J_{2n} deduced from the wheel W_{2n} by alternately deleting *n* spokes has unbounded metric dimension with $\beta(J_{2n}) = \lfloor \frac{2n}{3} \rfloor$ for every $n \ge 4$.

Recently, the metric dimension of some wheel related graphs and the convex polytopes generated by wheel related graph are studied by Imran et al. [14] and Siddiqui et al. [22], respectively, where they have provided an exact formula for the metric dimension of several class of wheel related graphs and shown that all those families have unbounded metric dimension. Further result about the unbounded metric dimension of graphs was derived by Abbas et al. in [1] and U. Ali et al. in [27]. In this context, it is natural to ask for characterization of graphs with unbounded metric dimension. In this paper, we have extended this study and determined a family of graph that is deduced from wheel graph by removing *k* consecutive spokes and shown that this family of graph has unbounded metric dimension. It is also worth mentioning that this family of graph have many interesting graph-theoretic properties and have been studied intensively in the literature; for example the chromaticity of this family was investigated in [7]. In this paper, the exact value of W(n, k) has also been studied, it is shown that exchange property does not hold for the resolving sets of graph W(n, k).

In what follows all indices *i* which do not satisfy inequalities $1 \le i \le n$ will be taken modulo *n*.

2. On metric basis of wheel graph with *k*-consecutive missing spokes

A wheel graph denoted by W_n is defined as $W_n \cong C_n + K_1$, Where $C_n : v_1, v_2, v_3, \ldots, v_n$ for $n \ge 3$ is cycle of length *n*. Suppose $v_1, v_2, v_3, \ldots, v_n$ are the vertices of the outer cycle C_n of W_n and *v* be the central vertex of W_n . If *p*, *q* be the positive integers such that $1 \le p < q \le n$, then $v_{p+1}, v_{p+2}, v_{p+3}, \ldots, v_{q-1}$ are the vertices in the *gap* determined by the vertices v_p and v_q , and the size of the gap is denoted by G_{q-p-1} is q - p - 1. Any two gaps G_{q-p-1} and G_{s-r-1} determined by the vertices v_p, v_q and v_r, v_s respectively are said to be the *neighboring gaps* if $v_q = v_r$.

Definition: A wheel graph with *k*-consecutive missing spokes denoted by W(n, k) can be obtained by deleting *k*-consecutive spokes from the wheel graph denoted by W_n . The graph W(n, k) is depicted in the Figure 1. In the following theorem, the metric dimension of the graph W(n, k) is determined. The proof of this theorem is given at the end of this section.



Figure 1. The wheel graph W(n, k) with *k*-consecutive missing spokes.

Theorem 2.1. Let W(n, k) denote the wheel graph with k-consecutive missing spokes with $k \ge 1$, then $\beta(W(n, k)) = \lceil \frac{2n-2k-2}{5} \rceil$ for $n \ge 5$.

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To prove this theorem, we need to prove some lemmas.

In the following result, the upper bound for the metric dimension of W(n, k) is derived.

Lemma 2.1. For every positive integers $n \ge 5$ and $k \ge 1$, we have

$$\beta(W(n,k) \le \lceil \frac{2n-2k-2}{5} \rceil.$$

Proof. Let

$$B = \{v_{\lceil \frac{5i}{2} \rceil} : 2 \le i \le \lceil \frac{2n - 2k - 2}{5} \rceil - 1\} \cup \{v_{n-k-1}, v_{n-\lfloor \frac{k-1}{2} \rfloor}\}$$

be a set of vertices of the graph W(n, k). In order to show that *B* is a resolving set for the graph W(n, k); we consider an arbitrary vertex $y \in V(W(n, k) \setminus B)$.

Then there are following possibilities for the choice of *y*:

• *The vertex y belong to a gap of size* $\lceil \frac{k}{2} \rceil + 3$ *of B.* From the construction of set *B*, it is evident that $\{v_1, v_2, v_3, v_4, v_{n-k+i} : \lceil \frac{k+1}{2} \rceil + 1 \le i \le k\}$ are the vertices in the gap of size $\lceil \frac{k}{2} \rceil + 3$ determined by the vertices $v_{n-k+\lceil \frac{k+1}{2} \rceil}$ and v_5 . The representation of these vertices are given below:

$$r(v_i|B) = \begin{cases} (k-i+3,...,k-i+3,i-\lceil\frac{k+1}{2}\rceil), & \lceil\frac{k+1}{2}\rceil+1 \le i \le k ;\\ (2,2,...,2,\frac{k+1}{2}), & i=1;\\ (2,2,...,2,\frac{k+1}{2}+1), & i=2;\\ (2,2,...,2,\frac{k+1}{2}+2), & i=3;\\ (2,2,...,2,\frac{k+1}{2}+2), & i=4. \end{cases}$$

Which shows that every vertex in the gap of size $\lceil \frac{k}{2} \rceil + 3$ has a unique representation (see Figure 2). • *y* belong to a gap of size $\lceil \frac{k}{2} \rceil$ of *B*. From the construction of set *B*, it is evident that $\{v_1, v_2, v_3, v_4, v_{n-k+i} : \lceil \frac{k+1}{2} \rceil + 1 \le i \le k\}$ are the vertices in the gap of size $\lceil \frac{k}{2} \rceil$ determined by the vertices $v_{n-k+\lceil \frac{k+1}{2} \rceil}$ and v_5 . The representation of these vertices are given below:

$$r(v_i|B) = \begin{cases} (i+2,i+2,...,i+2,i+1,\lceil\frac{k+1}{2}\rceil)-i, & 0 \le i \le \lceil\frac{k+1}{2}\rceil-1 ;\\ (2,2,...,2,\frac{k+1}{2}), & i = 1;\\ (2,2,...,2,\frac{k+1}{2}+1), & i = 2;\\ (2,2,...,2,\frac{k+1}{2}+2), & i = 3;\\ (2,2,...,2,\frac{k+1}{2}+2), & i = 4. \end{cases}$$

Which shows that every vertex in the gap of size $\lceil \frac{k}{2} \rceil + 3$ has a unique representation (see Figure 2).

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Figure 2. Resolving sets for graph W(20, 4) and W(20, 5).

• *y* belong to a gap of size 2 of *B*. Let v_{p+1} and v_{p+2} be the vertices in the gap of size 2 determined by the vertices v_p and v_{p+3} . If $y = v_{p+1}$, then it is the unique vertex that has distance 1 and 3 from v_p and v_n , respectively, and distance 2 from all other vertices of *B*. If $y = v_{p+2}$, then it is again the unique vertex that has distance 1 and 3 from v_3 and v_3 , respectively, and has distance 2 from all other vertices of *B*.

• *y* belong to a gap of size 1 of B. In this case, let v_{p+1} be the vertex in the gap of size one determined by the vertices v_p and v_{p+2} ; then the vertex v_{p+1} is the unique vertex that has distance 1 from v_p and v_{p+2} .

The above discussion implies that, every vertex of the graph W(n, k) has unique representation with respect to the set *B*. Hence *B* is a resolving set that show that

$$\beta(W(n,k) \le \lceil \frac{2n-2k-2}{5} \rceil$$

If we delete the vertices $v_{n-k-1}, v_{n-k-2}, \ldots, v_{n-1}, v_n$ from the graph denoted by W(n, k), then the vertex

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set of the new graph denoted by G_1 is

$$V(G_1) = V(W(n,k)) \setminus \{v_{n-k+i} : 1 \le i \le k\}.$$

From the construction, G_1 is isomorphic to the join of the path graph P_{n-k} and K_1 . That is $G_1 \cong P_{n-k} + K_1$, as shown in Figure 3.



Figure 3. The graph $G_1 \cong W(n,k) \setminus \{v_{n-k+i} : 1 \le i \le k\}$.

Lemma 2.2. Every basis of the graph G_1 satisfy the following conditions:

(i) A gap determine by any vertex from v_1 to v_{n-k} is at most of size 3.

(ii) There is at most one gap in the basis set of size 3.

(iii) At least one of v_{n-k} , v_{n-k-1} and v_{n-k-2} must belong to basis set.

(iv) The central vertex v_{\circ} does not belong to any of the basis.

(v) If a gap in any of basis contains at least two vertices; then its neighboring gap contain at most one vertex

Proof. Let B_1 be the arbitrary basis of graph denoted by G_1 . We prove that B_1 satisfy conditions (i) - (v).

(*i*): If there exist a gap of size 4 having the vertices $\{v_p, v_{p+1}, v_{p+2}, v_{p+3}\}$. But in this case, we have $r(v_{p+1}|B_1) = r(v_{p+2}|B_2) = (2, 2, 2, ..., 2, 2)$, a contradiction. Thus, any gap determine by any vertex from v_1 to v_{n-k} is at most of size 3.

(*ii*): Suppose on contrary, that there exist two gaps of size 3. Let $\{v_p, v_{p+1}, v_{p+2}\}$ and $\{v_q, v_{q+1}, v_{q+2}\}$ are the vertices in these two gaps. Then, we have $r(v_{p+1}|B) = r(v_{q+1}|B) = (2, 2, 2, ..., 2, 2)$, a contradiction. Hence, there is at most one gap in B_1 of size 3

(*iii*): At least one of v_{n-k} , v_{n-k-1} and v_{n-k-2} must belong to B_1 , otherwise v_{n-k} and v_{n-k-1} have same representation. Similarly one of the v_1 , v_2 and v_3 also belong to B_1 .

(*iv*): Since v_{\circ} is at the distance 1 to all the vertices of *G* and can be removed from any resolving set. Therefore B_1 being a minimal resolving set can not contain v_{\circ} . This implies that he central vertex v_{\circ} does not belong to B_1 .

(v): Let there are two consecutive gaps of size at least 2 and $\{v_{p+1}, v_{p+2}\}$ and $\{v_{p+4}, v_{p+5}\}$ be the vertices in these gaps determined by v_p , v_{p+3} and v_{p+3} , v_{p+6} , respectively. Then the vertices v_{p+2} and v_{p+4} have

same representations, a contradiction. Hence, if a gap of B_1 contains at least two vertices; then its neighboring gap contain at most one vertex.

Lemma 2.3. Every subset of the vertex set of the graph G_1 satisfying the conditions of Lemma 2.2 is a resolving set.

Proof. Let B_1 is a subset of $V(G_1$ satisfying all the conditions of Lemma 2.2. Let $y \in V(G_1 \setminus B_1)$ be an arbitrary element, we show that *y* has unique representation. There are following possibilities for the choice of *y*.

• *y* belong to a gap of size 3 of B_1 . Assume that the vertices $v_p, v_{p+1}, v_{p+2}, v_{p+3}$,

 v_{p+4} belong to the rim vertices; where v_p and $v_{p+4} \in B_1$. The vertex v_{p+1} then has distance one from v_p and 2 from all other vertices of B_1 , and v_{p+2} then has distance 2 from all vertices of B_1 . Similarly the vertex v_{p+3} is it distance 1 from v_{p+4} and has distance 2 from all other vertices of B_1 . By condition (*i*) and (*ii*) of Lemma 2.2; these vertices have unique representation.

• *y* belong to a gap of size 2 of B_1 . Let the vertices $v_p, v_{p+1}, v_{p+2}, v_{p+3}$ belong to the rim vertices such that v_p and $v_{p+3} \in B_1$. We have two cases here:

If $y_1 = v_{p+1}$, then it has distance 1 from v_p and distance 2 from all other vertices of B_1 . By condition (*v*) of Lemma 2.2; these vertices have unique representation.

If $y_1 = v_{p+2}$, then it is at distance 1 from v_{p+3} and distance 2 from all vertices of B_1 . Again, by condition (*v*) of Lemma 2.2; these vertices have unique representation and no other vertex have this representation.

• *y* belong to a gap of size 1 of B_1 . Assume that $\{v_p, y_1, v_{p+1}\}$ belong to the rim vertices; where v_p and $v_{p+1} \in B_1$. Then *y* is the only vertex having distance 1 from both the vertices v_p and v_{p+1} , so the representation of *y* is unique.

Thus, we conclude that any set B_1 satisfying all conditions of Lemma 2.2 is the resolving set of the graph $G_1 \cong P_{n-k} + K_1$.

 \Box In the following result, the metric dimension of the graph $G_1 \cong P_{n-k} + K_1$ is computed.

Lemma 2.4. For every positive integer $m \ge 3$, $\beta(G_1) = \lceil \frac{2m-2}{5} \rceil$, where m = n - k.

Proof. Define the set B_1 as follows,

$$B_1 = \{v_1\} \cup \{v_{\lceil \frac{5i-1}{2} \rceil}, 2 \le i \le \lceil \frac{2m-2}{5} \rceil - 1\} \cup \{v_{m-1}\}$$

The set B_1 satisfy all the conditions of Lemma 2.2. Therefore B_1 is the resolving set for G_1 . Which shows that

$$\beta(G_1) = \le |B_1| = \lceil \frac{2m-2}{5} \rceil.$$
(2.1)

To prove the lower bound, let B_1 be the basis for the graph G_1 . To show that $|B_1| \ge \lceil \frac{2m-2}{5} \rceil$, we have the following two cases:

Case 1: When $|B_1| = 2l$, where $l \ge 1$.

• If B_1 has a gap of size 3, then from Lemma 2.2, B_1 can contain at most one gap of size 3, l - 1 gaps of size 2, l gaps of size 1. So we have

$$m - B_1 \le 3 + 2(l - 1) + l$$

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$$\implies m - 2l \le 3 + 2l - 2 + l$$
$$\implies l \ge \frac{m - 1}{5}$$
$$\implies 2l = |B_1| \ge \frac{2m - 2}{5}$$
$$\implies |B_1| \ge \lceil \frac{2m - 2}{5} \rceil.$$

• If B_1 has no gap of size 3, then again from Lemma 2.2, B_1 contain no gap of size 3, at most l gaps of size 2 and l gaps of size 1. Hence we have

$$m - B_1 \le 2l + l$$
$$\implies m - 2l \le 3l \implies l \ge \frac{m}{5}$$
$$\implies 2l = |B_1| \ge \frac{2m}{5}$$
$$\implies |B_1| \ge \lceil \frac{2m}{5} \rceil \ge \lceil \frac{2m - 2}{5} \rceil.$$

Case 2: When $|B_1| = 2l + 1$, where $l \ge 1$.

• If B_1 has a gap of size 3; then from Lemma 2.2, B_1 can contain at most one gap of size 3, l - 1 gaps of size 2 and l gaps of size 1. Therefore, we have

$$m - B_1 \leq 3 + 2(l - 1) + l$$

$$\implies m - (2l + 1) \leq 3 + 2l - 2 + l$$

$$\implies l \geq \frac{m - 2}{5} \implies 2l + 1 \geq \frac{2(m - 2)}{5} + 1$$

$$\implies 2l + 1 = |B_1| \geq \frac{2m + 1}{5}$$

$$\implies |B_1| \geq \lceil \frac{2m + 1}{5} \rceil \geq \lceil \frac{2m - 2}{5} \rceil.$$

• If B_1 has no gap of size 3, then again from Lemma 2.2, B_1 contain no gap of size 3, at most l gaps of size 2 and l gaps of size 1. We have

$$m - B_1 \le 2l + l$$

$$\implies m - (2l + 1) \le 2l + l$$

$$\implies l \ge \frac{m - 1}{5} \implies 2l + 1 \ge \frac{2(m - 1)}{5} + 1$$

$$\implies 2l + 1 = |B_1| \ge \frac{2m + 3}{5}$$

$$\implies |B_1| \ge \frac{2m + 3}{5}$$

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$$\implies |B_1| \ge \lceil \frac{2m+3}{5} \rceil \ge \lceil \frac{2m-2}{5} \rceil.$$

The above discussion implies that

 $\beta(G_1) \ge \lceil \frac{2m-2}{5} \rceil \tag{2.2}$

From Eqs (2.1) and (2.2), we get

$$\beta(G_1) = \lceil \frac{2m-2}{5} \rceil.$$

We now are in position to prove the main theorem of our paper.

Proof of Theorem 2.1: The upper bound for the metric dimension of the wheel graph with k-consecutive missing spokes denoted by W(n, k) is given by Lemma 2.1:

$$\beta(G_1) = \beta(W(n,k)) \le \lceil \frac{2n-2k-2}{5} \rceil, \tag{2.3}$$

where $n \ge 5$ and $k \ge 1$. On the other hand from the Lemma 2.4, we have

$$\beta(G_1) \le \lceil \frac{2m-2}{5} \rceil,$$

where $G_1 \cong W(n,k) \setminus \{v_{n-k+i} : 1 \le i \le k\}$. Now by putting the value of m = n - k, we get

$$\beta(G_1) \ge \lceil \frac{2m-2}{5} \rceil = \lceil \frac{2n-2k-2}{5} \rceil.$$

Since G_1 is the subgraph of graph W(n, k), therefore

$$\beta(W(n,k)) \ge \lceil \frac{2n-2k-2}{5} \rceil.$$
(2.4)

Hence, from Equations (3) and (4), it follows that

$$\beta(W(n,k)) = \lceil \frac{2n-2k-2}{5} \rceil,$$

where $n \ge 5$ and $k \ge 1$. This completes the proof.

3. Exchange property for resolving set in W(n, k)

We have seen that a subset *W* of vertices of a graph *G* is a resolving set if every vertex in *G* is uniquely determined by its distances to the vertices of *W*. Resolving sets behave like bases in a *vector space* in that each vertex in the graph can be uniquely identified relative to the vertices of these sets. But though resolving sets do share some of the properties of bases in a vector space, they do not always have the exchange property from linear algebra. Resolving sets are said to have the *exchange property* in *G* if whenever *S* and *R* are minimal resolving sets for *G* and $r \in R$, then there exists $s \in S$ so that $(S \setminus \{s\}) \cup \{r\}$ is a minimal resolving set [2].

If the exchange property holds for a graph G, then every minimal resolving set for G has the same size and algorithmic methods for finding the metric dimension of G are more feasible. Thus to show

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that the exchange property does not hold in a given graph, it is sufficient to show two minimal resolving sets of different size. However, since the converse is not true, knowing that the exchange property does not hold does not guarantee that there are minimal resolving sets of different size.

The following results concerning exchange property for resolving sets were deduced in [2].

Theorem 3.1. [2] The exchange property holds for resolving sets in trees.

Theorem 3.2. [2] For $n \ge 8$, resolving sets do not have the exchange property in wheels W_n .

The exchange property for resolving sets of the graph W(n, k) has been discussed in the next theorem.

Theorem 3.3. The exchange property does not hold for the resolving sets of the graph W(n, k) for every positive integers $n - k \ge 10$ and $k \ge 1$.

Proof. To show that exchange property does not hold on the resolving sets of graph W(n, k), we need to show that there are two minimal resolving sets of different size.

Since the set $B = \{v_{\lceil \frac{5i}{2} \rceil} : 2 \le i \le \lceil \frac{2n-2k-2}{5} \rceil - 1\} \cup \{v_{n-k-1}, v_{n-\lfloor \frac{k-1}{2} \rfloor}\}$ is the metric basis of the graph W(n, k), so indeed it is a minimal resolving set for W(n, k) having cardinality m = n - k.

Now define another subset of the vertex set of the graph W(n, k), for k + 10 and $k \ge 1$ as follows:

$$B^* = \{v_{n-\lfloor \frac{k}{2} \rfloor}, v_{4i+1}, v_{4j+2} : 1 \le i \le \lceil \frac{n-2k-2}{4} \rceil : 1 \le j \le \lceil \frac{n-2k}{4} \rceil - 1\} \\ \cup \{v_{n-k}\} \cup \{v_{n-\lfloor \frac{k-1}{2} \rfloor}\}$$

 B^* satisfy all the condition of Lemma 2.2, so is a resolving set of W(n,k) having cardinality $\lceil \frac{n-2k-2}{4} \rceil + \lceil \frac{n-2k}{4} \rceil + 1$. Next, we show that B^* is also minimal resolving set by proving that there does not exist any $b \in B^*$ such that $B^* \setminus \{b\}$ is still a resolving set. There are the following three cases for the choice of *b*:

Case 1: If $b = v_{n-\lfloor \frac{k-1}{2} \rfloor}$; then removal of $v_{n-\lfloor \frac{k}{2} \rfloor}$ would yield a gap of size k + 2. By claim (1), this is not possible.

Case 2: If $b \in \{v_{4i+1} : 1 \le i \le \lceil \frac{n-2k-2}{4} \rceil\}$ then removal of *b* can be considered as follows: If $b = v_5$ for i = 1, then it would cause two vertices having same representation as $r(v_5|B^*) = r(v_7|B^*) = (1, 2, 2, 2, ..., \lfloor \frac{k}{2} \rfloor + 3)$. When $b = v_9$ for i = 1, then it would cause two vertices having same representation, i.e. $r(v_9|B^*) = r(v_{11}|B^*) = (2, 2, 1, 2, ..., \lfloor \frac{k}{2} \rfloor + 3)$. Similarly, we can show that if we remove any basis vertex v_{4i+1} for $1 \le i \le \lceil \frac{n-2k-2}{4} \rceil$, then v_{4i+1} and v_{4i+3} will have the same representations.

any basis vertex v_{4i+1} for $1 \le i \le \lceil \frac{n-2k-2}{4} \rceil$, then v_{4i+1} and v_{4i+3} will have the same representations. **Case 3**: If $b \in \{v_{4j+2} : 1 \le j \le \lceil \frac{n-2k}{4} \rceil - 1\}$. In this case by removing $b = v_6$ for j = 1, would cause two vertices having same representations which are $r(v_5|B^*) = r(v_7|B^*) = (1, 2, 2, 2, ..., \lfloor \frac{k-1}{2} \rfloor + 3)$. If $b = v_9$ for i = 1, then it would cause two vertices having same representations, i.e. $r(v_9|B^*) = r(v_{11}|B^*) = (2, 2, 1, 2..., \lfloor \frac{k-1}{2} + 3 \rfloor)$. Similarly, it can be shown that for any choice of the vertex *b* from the set of vertices $\{v_{4j+2} : 1 \le j \le \lceil \frac{2n-2k}{5} \rceil - 1\}$; the vertices v_{4j+1} and v_{4j+3} will have the same representations. This shows that B^* is a minimal resolving set, which completes the proof.

4. Conclusions

The problem of determining whether $\beta(G) < k$ is an *NP*-complete problem. It is natural to ask for characterization of graphs with bounded or unbounded metric dimension. In general, it appears

unlikely that significant progress can be made in this regard unless it belongs to a class for which the distances between vertices can be described in some systematic manner. There have been significant work done in literature to explore the families of graphs with unbounded metric dimension (See [1, 3, 14, 22, 26, 27]). In this context, we have explore a family that is constructed from wheel graph by removing k consecutive spokes and it is shown that this family of graph has unbounded metric dimension. This results also add further support to a negative answer of an open problem raised in [12]. It is further shown that exchange property for resolving set does not hold for this family of graphs.

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Conflict of interest

The authors declare no conflict of interest.

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